

Why Functions as Values

- Abstraction is easier with functions as values
 - `get-ids`
 - `filter`, `map`, etc.
- Separate `defun` form becomes unnecessary
 - ```
{defun {f x} {+ 1 x}}
{f 10}
```

⇒

```
{with {f {fun {x} {+ 1 x}}}
 {f 10}}
```

# FWAE Grammar, Almost

```
<FWAE> ::= <num>
| { + <FWAE> <FWAE> }
| { - <FWAE> <FWAE> }
| { with { <id> <FWAE> } <FWAE> }
| <id>
| { <id> <FWAE> }
| { fun { <id> } <FWAE> }
```

?



# FWAE Evaluation

10  $\Rightarrow$  10

{+ 1 2}  $\Rightarrow$  3

{- 1 2}  $\Rightarrow$  -1

{with {x 7} {+ x 2}}  $\Rightarrow$  {+ 7 2}  $\Rightarrow$  9

y  $\Rightarrow$  *free variable*

{fun {x} {+ 1 x}}  $\Rightarrow$   
{fun {x} {+ 1 x}}

Result is not always a number!

; interp FWAE ... -> FWAE-Value

# FWAE Evaluation

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{with {x 7} {+ x 2}}  $\Rightarrow$  {+ 7 2}  $\Rightarrow$  9

y  $\Rightarrow$  *free variable*

{fun {x} {+ 1 x}}  $\Rightarrow$   
{fun {x} {+ 1 x}}

{with {y 10} {fun {x} {+ y x}}}  
 $\Rightarrow$  {fun {x} {+ 10 x}}

{with {f {fun {x} {+ 1 x}}}} {f 3}}  
 $\Rightarrow$  {{fun {x} {+ 1 x}} 3}

Doesn't match the grammar for <FWAE>

# FWAE Grammar

```
<FWAE> ::= <num>
| { + <FWAE> <FWAE> }
| { - <FWAE> <FWAE> }
| { with { <id> <FWAE> } <FWAE> }
| <id>
| { <id> <FWAE> }
| { fun { <id> } <FWAE> }
| { <FWAE> <FWAE> }
```

NEW

NEW

# FWAE Evaluation

$\{\text{with } \{\text{f } \{\text{fun } \{\text{x}\} \{+ 1 \text{x}\}\}\} \{\text{f } 3\}\}$   
 $\Rightarrow \{\{\text{fun } \{\text{x}\} \{+ 1 \text{x}\}\} 3\}$   
 $\Rightarrow \{+ 1 3\} \Rightarrow 4$

$\{\{\text{fun } \{\text{x}\} \{+ 1 \text{x}\}\} 3\} \Rightarrow \{+ 1 3\} \Rightarrow 4$

$\{1 2\} \Rightarrow \textit{not a function}$

$\{+ 1 \{\text{fun } \{\text{x}\} 10\}\} \Rightarrow \textit{not a number}$

# FWAE Datatype

```
(define-type FWAE
 [num (n number?)]
 [add (lhs FWAE?)
 (rhs FWAE?)]
 [sub (lhs FWAE?)
 (rhs FWAE?)]
 [with (name symbol?)
 (named-expr FWAE?)
 (body FWAE?)]
 [id (name symbol?)]
 [fun (param symbol?)
 (body FWAE?)]
 [app (fun-expr FWAE?)
 (arg-expr FWAE?)])
```

```
(test (parse '{fun {x} {+ x 1}})
 (fun 'x (add (id 'x) (num 1))))
```

# FWAE Datatype

```
(define-type FWAE
 [num (n number?)]
 [add (lhs FWAE?)
 (rhs FWAE?)]
 [sub (lhs FWAE?)
 (rhs FWAE?)]
 [with (name symbol?)
 (named-expr FWAE?)
 (body FWAE?)]
 [id (name symbol?)]
 [fun (param symbol?)
 (body FWAE?)]
 [app (fun-expr FWAE?)
 (arg-expr FWAE?)])
```

```
(test (parse '{fun {x} {+ x 1}} 10)
 (app (fun 'x (add (id 'x) (num 1))) (num 10)))
```



# FWAE Interpreter

```
; interp : FWAE -> FWAE
(define (interp a-wae)
 (type-case FWAE a-wae
 [num (n) a-wae]
 [add (l r) (num+ (interp l) (interp r))]
 [sub (l r) (num- (interp l) (interp r))]
 [with (bound-id named-expr body-expr)
 (interp (subst body-expr
 bound-id
 (interp named-expr)))]
 [id (name) (error 'interp "free variable")]
 [fun (param body-expr)
 a-wae]
 [app (fun-expr arg-expr)
 (local [(define fun-val (interp fun-expr))]
 (interp (subst (fun-body fun-val)
 (fun-param fun-val)
 (interp arg-expr))))))])
```

# Add and Subtract

```
; num+ : FWAE FWAE -> FWAE
(define (num+ x y)
 (num (+ (num-n x) (num-n y))))
```

```
; num- : FWAE FWAE -> FWAE
(define (num- x y)
 (num (- (num-n x) (num-n y))))
```

Better:

```
; num-op : (num num -> num) -> (FWAE FWAE -> FWAE)
(define (num-op op)
 (lambda (x y)
 (num (op (num-n x) (num-n y)))))
```

```
(define num+ (num-op +))
(define num- (num-op -))
```

# FWAE Subst

**; subst : FWAE symbol -> FWAE**

Implementation is an exercise for the student

Beware: with the obvious implementation,

```
(subst {with {y 10} z} 'z {fun {x} {+ x y}})
⇒ {with {y 10} {fun {x} {+ x y}}}
```

which is wrong, but we leave this problem to CS  
7520

- Only happens when the original program has free variables
- The problem disappears with deferred substitution, anyway

# No More With

Compare the `with` and `app` implementations:

```
(define (interp a-wae)
 (type-case FWAE a-wae
 ...
 [with (bound-id named-expr body-expr)
 (interp (subst body-expr
 bound-id
 (interp named-expr)))]
 ...
 [app (fun-expr arg-expr)
 (local [(define fun-val (interp fun-expr))]
 (interp (subst (fun-body fun-val)
 (fun-param fun-val)
 (interp arg-expr))))]))
```

The `app` case does everything that `with` does

# No More With

```
{with {x 10} x}
```

is the same as

```
{{fun {x} x} 10}
```

In general,

```
{with {<id> <FWAE>1} <FWAE>2}
```

is the same as

```
{{fun {<id>} <FWAE>2} <FWAE>1}
```

Let's assume

```
(test {with {<id> <FWAE>1} <FWAE>2}
 (app (fun {<id> <FWAE>2}
 <FWAE>1)))
```

# FAE Grammar

```
<FAE> ::= <num>
 | { + <FAE> <FAE> }
 | { - <FAE> <FAE> }
 | { with { <id> <FAE> } <FAE> }
 | <id>
 | { fun { <id> } <FAE> }
 | { <FAE> <FAE> }
```

- We'll still use `with` in boxes
- No more case lines in `interp`, etc. for `with`
- No more test cases for `interp`, etc. using `with`

# FAE Interpreter

```
; interp : FAE -> FAE
(define (interp a-fae)
 (type-case FAE a-wae
 [num (n) a-fae]
 [add (l r) (num+ (interp l) (interp r))]
 [sub (l r) (num- (interp l) (interp r))]
 [id (name) (error 'interp "free variable")]
 [fun (param body-expr) a-fae]
 [app (fun-expr arg-expr)
 (local [(define fun-val (interp fun-expr))]
 (interp (subst (fun-body fun-val)
 (fun-param fun-val)
 (interp arg-expr))))]))))
```

# FAE with Deferred Substitution

(interp {with {y 10} {fun {x} {+ y x}}})

⇒

(interp {fun {x} {+ y x}})

(interp {{fun {y} {fun {x} {+ y x}}} 10})

⇒

(interp {fun {x} {+ y x}})



# F AE with Deferred Substitution

```
(interp {{with {y 10} {fun {x} {+ y x}}}}
 {with {y 7} y}})
```

Argument expression:

```
(interp {with {y 7} y})
```

⇒

```
(interp y) ⇒ 7
```

Function expression:

```
(interp {with {y 10} {fun {x} {+ y x}}}})
```

⇒

```
(interp {fun {x} {+ y x}}) ⇒ ?
```

# F&A;E Values

A function value needs to keep its substitution cache

```
(define-type F&A;E-Value
 [numV (n number?)]
 [closureV (param symbol?)
 (body F&A;E?)
 (ds DefrdSub?)])
```

```
(define-type DefrdSub
 [mtSub]
 [aSub (name symbol?)
 (value F&A;E-Value?)
 (ds DefrdSub?)])
```

```
(test (interp {with {y 10} {fun {x} {+ y x}}})
 (closureV 'x {+ y x}
 (aSub 'y (num 10) (mtSub))))
```

# Continuing Evaluation

Function: `{fun {x} {+ y x}}` y = 10

Argument: 7

To apply, interpret the function body with the given argument:

`(interp {+ y x})` x = 7    y = 10

# F AE Interpreter with Substitution

```
; interp : FAE DefrdSub -> FAE-Value
(define (interp a-wae ds)
 (type-case FAE a-wae
 [num (n) (numV n)]
 [add (l r) (num+ (interp l ds) (interp r ds))]
 [sub (l r) (num- (interp l ds) (interp r ds))]
 [id (name) (lookup name ds)]
 [fun (param body-expr)
 (closureV param body-expr ds)]
 [app (fun-expr arg-expr)
 (local [(define fun-val
 (interp fun-expr ds))
 (define arg-val
 (interp arg-expr ds))]
 (interp (closureV-body fun-val)
 (aSub (closureV-param fun-val)
 arg-val
 (closureV-ds fun-val))))))])
```