

Arithmetic Language

<AE> ::= <num>
| {+ <AE> <AE>}
| {- <AE> <AE>}

```
(define-type AE  
  [num (n number?) ]  
  [add (lhs AE?)  
    (rhs AE?) ]  
  [sub (lhs AE?)  
    (rhs AE?) ] )
```

Arithmetic Language

```
<AE> ::= <num>
      | {+ <AE> <AE>}
      | {- <AE> <AE>}
```

```
; interp AE -> num
(define (interp an-ae)
  (type-case AE an-ae
    [num (n) n]
    [add (l r) (+ (interp l) (interp r))]
    [sub (l r) (- (interp l) (interp r))]))
```

```
(test (interp (add (num 1) (add 2))) 3)
```

Arithmetic Language

```
<AE> ::= <num>
      | {+ <AE> <AE>}
      | {- <AE> <AE>}
```

```
; parse: s-expr -> AE
(define (parse input)
  (cond
    [(number? input) (num input)]
    [(eq? (first input) '+)
     (add (parse (second input))
          (parse (third input)))]
    [(eq? (first input) '-')
     (sub (parse (second input))
          (parse (third input)))]
    [else (error 'parse "bad syntax: ~a" input)]))

(test (interp (parse '{+ 1 2})) 3)
```

The Arbitrariness of Identifiers

Are the following two programs equivalent?

```
(define (f x) (+ x 1))  
(f 10)
```

```
(define (f y) (+ y 1))  
(f 10)
```

yes

argument is consistently renamed

The Arbitrariness of Identifiers

Are the following two programs equivalent?

```
(define (f x) (+ x 1))  
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```
(define (f x) (+ y 1))  
(f 10)
```

no

not a use of the argument anymore

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```
(define (f y) (+ x 1))  
(f 10)
```

no

not a use of the argument anymore

The Arbitrariness of Identifiers

Are the following two programs equivalent?

```
(define (f x) (+ y 1))  
(f 10)
```

```
(define (f z) (+ y 1))  
(f 10)
```

yes

argument never used, so almost any name is ok

The Arbitrariness of Identifiers

Are the following two programs equivalent?

```
(define (f x) (+ y 1))  
(f 10)
```

```
(define (f y) (+ y 1))  
(f 10)
```

no

now a use of the argument

The Arbitrariness of Identifiers

Are the following two programs equivalent?

```
(define (f x) (+ y 1))  
(f 10)
```

```
(define (f x) (+ z 1))  
(f 10)
```

no

still an undefined identifier, but a different one

The Arbitrariness of Identifiers

Are the following two programs equivalent?

```
(define (f x)
  (local [(define y 10)]
    (+ x y)))
(f 0)
```

```
(define (f z)
  (local [(define y 10)]
    (+ z y)))
(f 0)
```

yes

argument is consistently renamed

The Arbitrariness of Identifiers

Are the following two programs equivalent?

```
(define (f x)
  (local [(define y 10)]
    (+ x y)))
(f 0)
```

```
(define (f x)
  (local [(define z 10)]
    (+ x z)))
(f 0)
```

yes

local identifier is consistently renamed

The Arbitrariness of Identifiers

Are the following two programs equivalent?

```
(define (f x)
  (local [(define y 10)]
    (+ x y)))
(f 0)
```

```
(define (f x)
  (local [(define x 10)]
    (+ x x)))
(f 0)
```

no

local identifier now hides the argument

The Arbitrariness of Identifiers

Are the following two programs equivalent?

```
(define (f x)
  (local [(define y 10)]
    (+ x y)))
(f 0)
```

```
(define (f y)
  (local [(define y 10)]
    (+ y y)))
(f 0)
```

no

local identifier now hides the argument

Free and Bound Identifiers

An identifier for the argument of a function or the name of a local identifier is a ***binding occurrence***

```
(define (f x y) (+ x y z))
```

```
(local [(define a 3)
        (define c 4)]
  (+ a b c))
```

Free and Bound Identifiers

A use of a function argument or a local identifier is a ***bound occurrence***

```
(define (f x y) (+ x y z))
```

```
(local [(define a 3)
        (define c 4)]
  (+ a b c))
```

Free and Bound Identifiers

A use of an identifier that is not function argument or a local identifier is a ***free identifier***

```
(define (f x y) (+ x y z))
```

```
(local [(define a 3)
        (define c 4)]
  (+ a b c))
```




Arithmetic Language

$\langle \text{AE} \rangle ::= \langle \text{num} \rangle$
| $\{ + \langle \text{AE} \rangle \langle \text{AE} \rangle \}$
| $\{ - \langle \text{AE} \rangle \langle \text{AE} \rangle \}$

No identifiers to help us study binding...

With Arithmetic Language

```
<WAE> ::= <num>
        | {+ <WAE> <WAE>}
        | {- <WAE> <WAE>}
        | {with {<id> <WAE>} <WAE>}
        | <id>
```

```
{with {x {+ 1 2}}
      {+ x x}} ⇒ 6
```

With Arithmetic Language

```
<WAE> ::= <num>
        | {+ <WAE> <WAE>}
        | {- <WAE> <WAE>}
        | {with {<id> <WAE>} <WAE>}
        | <id>
```

NEW

NEW

x \Rightarrow *error: free identifier*

With Arithmetic Language

```
<WAE> ::= <num>
        | {+ <WAE> <WAE>}
        | {- <WAE> <WAE>}
        | {with {<id> <WAE>} <WAE>}
        | <id>
```

NEW

NEW

```
{+ {with {x {+ 1 2}}
      {+ x x}}
  {with {x {- 4 3}}
      {+ x x}}}} ⇒ 8
```

With Arithmetic Language

```
<WAE> ::= <num>
        | {+ <WAE> <WAE>}
        | {- <WAE> <WAE>}
        | {with {<id> <WAE>} <WAE>}
        | <id>
```





```
{+ {with {x {+ 1 2}}
      {+ x x}}
  {with {y {- 4 3}}
      {+ y y}}}
```

⇒ 8

With Arithmetic Language



```
<WAE> ::= <num>
        | {+ <WAE> <WAE>}
        | {- <WAE> <WAE>}
        | {with {<id> <WAE>} <WAE>}
        | <id>
```

```
{with {x {+ 1 2}}
  {with {x {- 4 3}}
    {+ x x}}}} ⇒ 2
```

With Arithmetic Language

```
<WAE> ::= <num>
        | {+ <WAE> <WAE>}
        | {- <WAE> <WAE>}
        | {with {<id> <WAE>} <WAE>}
        | <id>
```



 

```
{with {x {+ 1 2}}
  {with {y {- 4 3}}
    {+ x x}}}
```

\Rightarrow 6

With Arithmetic Language



```
<WAE> ::= <num>
        | {+ <WAE> <WAE>}
        | {- <WAE> <WAE>}
        | {with {<id> <WAE>} <WAE>}
        | <id>
```

```
(define-type WAE
  [num (n number?)]
  [add (lhs WAE?)
       (rhs WAE?)]
  [sub (lhs WAE?)
       (rhs WAE?)]
  [with (name symbol?)
        (named-expr WAE?)
        (body WAE?)]
  [id (name symbol?)])
```


With Arithmetic Language



```
<WAE> ::= <num>
        | {+ <WAE> <WAE>}
        | {- <WAE> <WAE>}
        | {with {<id> <WAE>} <WAE>}
        | <id>
```

```
(define (interp a-wae)
  (type-case WAE a-wae
    [num (n) n]
    [add (l r) (+ (interp l) (interp r))]
    [sub (l r) (- (interp l) (interp r))]
    [with (bound-id named-expr body-expr)
     ...]
    [id (name) ...])))
```

With Arithmetic Language


```
<WAE> ::= <num>
        | {+ <WAE> <WAE>}
        | {- <WAE> <WAE>}
        | {with {<id> <WAE>} <WAE>}
        | <id>
```

```
(define (interp a-wae)
  (type-case WAE a-wae
    [num (n) n]
    [add (l r) (+ (interp l) (interp r))]
    [sub (l r) (- (interp l) (interp r))]
    [with (bound-id named-expr body-expr)
     ...]
    [id (name) (error 'interp "free variable")]))
```

With Arithmetic Language



```
<WAE> ::= <num>
        | {+ <WAE> <WAE>}
        | {- <WAE> <WAE>}
        | {with {<id> <WAE>} <WAE>}
        | <id>
```



```
(define (interp a-wae)
  (type-case WAE a-wae
    [num (n) n]
    [add (l r) (+ (interp l) (interp r))]
    [sub (l r) (- (interp l) (interp r))]
    [with (bound-id named-expr body-expr)
      ... (interp named-expr) ... ]
    [id (name) (error 'interp "free variable")]))
```

With Arithmetic Language



```
<WAE> ::= <num>
        | {+ <WAE> <WAE>}
        | {- <WAE> <WAE>}
        | {with {<id> <WAE>} <WAE>}
        | <id>
```

```
(define (interp a-wae)
  (type-case WAE a-wae
    [num (n) n]
    [add (l r) (+ (interp l) (interp r))]
    [sub (l r) (- (interp l) (interp r))]
    [with (bound-id named-expr body-expr)
     ... (interp named-expr)
     ... (interp body-expr) ... ]
    [id (name) (error 'interp "free variable")]))
```

With Arithmetic Language

```
<WAE> ::= <num>
        | {+ <WAE> <WAE>}
        | {- <WAE> <WAE>}
        | {with {<id> <WAE>} <WAE>}
        | <id>
```

```
(define (interp a-wae)
  (type-case WAE a-wae
    [num (n) n]
    [add (l r) (+ (interp l) (interp r))]
    [sub (l r) (- (interp l) (interp r))]
    [with (bound-id named-expr body-expr)
      (interp (subst body-expr
                     bound-id
                     (interp named-expr)))]
    [id (name) (error 'interp "free variable")]))
```

Substitution

```
; subst : WAE symbol num -> WAE  
(define (subst a-wae sub-id val)  
  (type-case WAE a-wae  
    [num (n) ...]  
    [add (l r) ...]  
    [sub (l r) ...]  
    [with (bound-id named-expr body-expr)  
      ...]  
    [id (name) ...]))
```

Let's make examples/tests first...

Example Substitutions

```
; 10 for x in {+ 1 x} ⇒ {+ 1 10}
(test (subst (add (num 1) (id 'x)) 'x 10)
      (add (num 1) (num 10)))
```

```
; 10 for x in x ⇒ 10
(test (subst (id 'x) 'x 10)
      (num 10))
```

```
; 10 for x in y ⇒ y
(test (subst (id 'y) 'x 10)
      (id 'y))
```

```
; 10 for y in {- x 1} ⇒ {- x 1}
(test (subst (sub (id 'x) (num 1)) 'y 10)
      (sub (id 'x) (num 1)))
```

Substitution

```
; subst : WAE symbol num -> WAE
(define (subst a-wae sub-id val)
  (type-case WAE a-wae
    [num (n) a-wae]
    [add (l r) (add (subst l sub-id val)
                   (subst r sub-id val))]
    [sub (l r) (sub (subst l sub-id val)
                   (subst r sub-id val))]
    [with (bound-id named-expr body-expr)
      ...]
    [id (name) (if (symbol=? name sub-id)
                   (num val)
                   a-wae))]))
```


Example Substitutions

```
; 10 for x in {with {y 17} x} ⇒ {with {y 17} 10}
(test (subst (with 'y (num 17) (id 'x)) 'x 10)
      (with 'y (num 17) (num 10)))
```

```
; 10 for x in {with {y x} y} ⇒ {with {y 10} y}
(test (subst (with 'y (id 'x) (id 'y)) 'x 10)
      (with 'y (num 10) (id 'y)))
```

```
; 10 for x in {with {x y} x} ⇒ {with {x y} x}
(test (subst (with 'x (id 'y) (id 'x)) 'x 10)
      (with 'x (id 'y) (id 'x)))
```

Substitution replaces

- free identifiers with the same name
- no binding identifiers
- no bound identifiers

Substitution

An identifier is bound when it appears in the body of a **with** binding the same name

Conversely, a free variable of a name appears in a **with** only if the **with** doesn't bind the name

```
; subst : WAE symbol num -> WAE  
(define (subst a-wae sub-id val)  
  (type-case WAE a-wae  
    ...  
    [with (bound-id named-expr body-expr)  
      (with bound-id  
        (subst named-expr sub-id val)  
        (if (symbol=? bound-id sub-id)  
          body-expr  
          (subst body-expr sub-id val))))]  
    ...))
```