

Recursion

```
{with {mk-rec {fun {body}
              {{fun {fX} {fX fX}}
               {fun {fX}
                 {{fun {f} {body f}}
                  {fun {x} {{fX fX} x}}}}}}}}
{with {fib {mk-rec
          {fun {fib}
            {fun {n}
              {if0 n
                 1
                 {if0 {- n 1}
                    1
                    {+ {fib {- n 1}}
                       {fib {- n 2}}}}}}}}}}
      {fib 4}}}
```

Typed Recursion

```
{with {mk-rec : (((num -> num) -> (num -> num)) -> (num -> num))
  {fun {body : ((num -> num) -> (num -> num))}
    {{fun {fX : ... -> (num -> num)} {fX fX}}
     {fun {fX : ... -> (num -> num)}
       {{fun {f : (num -> num)} {body f}}
        {fun {x : num} {{fX fX} x}}}}}}}}
{with {fib : (num -> num)}
  {mk-rec
    {fun {fib : (num -> num)}
      {fun {n : num}
        {if0 n
          1
          {if0 {- n 1}
            1
            {+ {fib {- n 1}}
              {fib {- n 2}}}}}}}}}}
  {fib 4}}}
```

Nothing works in place of ...

Extending the Type System

When the type system rejects your perfectly good program, it may be time to extend the type system

In this case, we can add **rec** as a core form, again

```
{rec {fib : (num -> num)
      {fun {n : num}
          {if0 n
              1
              {if0 {- n 1}
                  1
                  {+ {fib {- n 1}}
                    {fib {- n 2}}}}}}}}}}
      {fib 4}}
```

TRCFAE Grammar

```
<TRCFAE> ::= <num>
           | {+ <TRCFAE> <TRCFAE>}
           | {- <TRCFAE> <TRCFAE>}
           | <id>
           | {fun {<id> : <TE>} <TRCFAE>}
           | {<TRCFAE> <TRCFAE>}
           | {if0 <TRCFAE> <TRCFAE> <TRCFAE>}
           | {rec {<id> : <TE> <TRCFAE>} <TRCFAE>}
<TE> ::= num
       | (<TE> -> <TE>)
```

NEW

TRCFAE Datatypes

```
(define-type FAE
  ...
  [rec (name : symbol)
       (ty : TE)
       (rhs-expr : FAE)
       (body-expr : FAE)])
```

TRCFAE Interpreter

```
(define (interp a-fae ds)
  (type-case FAE a-fae
    ...
    [rec (bound-id type named-expr body-expr)
      (local [(define value-holder (box (numV 42)))
              (define new-ds (aRecSub bound-id
                                      value-holder
                                      ds))])
      (begin
        (set-box! value-holder (interp named-expr new-ds))
        (interp body-expr new-ds))))))
```

TRCFAE Interpreter Lookup

```
(define (lookup name ds)
  (type-case DefrdSub ds
    [mtSub () (error 'lookup "free variable")]
    [aSub (sub-name val rest-ds)
      (if (symbol=? sub-name name)
          val
          (lookup name rest-ds))]
    [aRecSub (sub-name val-box rest-ds)
      (if (symbol=? sub-name name)
          (unbox val-box)
          (lookup name rest-ds))]))
```

TRCFAE Type Checker

```
(define typecheck : (FAE TypeEnv -> Type)
  (lambda (fae env)
    (type-case FAE fae
      ...
      [rec (name ty rhs-expr body-expr)
         ....])))
```

$$\frac{\Gamma[\langle \text{id} \rangle \leftarrow \tau_0] \vdash \mathbf{e}_0 : \tau_0 \quad \Gamma[\langle \text{id} \rangle \leftarrow \tau_0] \vdash \mathbf{e}_1 : \tau_1}{\Gamma \vdash \{\text{rec } \{\langle \text{id} \rangle : \tau_0 \ \mathbf{e}_0\} \ \mathbf{e}_1\} : \tau_1}$$

TRCFAE Type Checker

```
(define typecheck : (FAE TypeEnv -> Type)
  (lambda (fae env)
    (type-case FAE fae
      ...
      [rec (name ty rhs-expr body-expr)
        (local [(define rhs-ty (parse-type ty))]
          ....
          (typecheck rhs-expr ....)
          (typecheck body-expr ....)
          ....)])))))
```

$$\frac{\Gamma[\langle id \rangle \leftarrow \tau_0] \vdash \mathbf{e}_0 : \tau_0 \quad \Gamma[\langle id \rangle \leftarrow \tau_0] \vdash \mathbf{e}_1 : \tau_1}{\Gamma \vdash \{\mathbf{rec} \{\langle id \rangle : \tau_0 \ \mathbf{e}_0\} \ \mathbf{e}_1\} : \tau_1}$$

TRCFAE Type Checker

```
(define typecheck : (FAE TypeEnv -> Type)
  (lambda (fae env)
    (type-case FAE fae
      ...
      [rec (name ty rhs-expr body-expr)
        (local [(define rhs-ty (parse-type ty))
                (define new-ds (aBind name
                                      rhs-ty
                                      env))])
          ....
          (typecheck rhs-expr new-ds)
          (typecheck body-expr new-ds)
          ....)]))
```

$$\frac{\Gamma[\langle \text{id} \rangle \leftarrow \tau_0] \vdash \mathbf{e}_0 : \tau_0 \quad \Gamma[\langle \text{id} \rangle \leftarrow \tau_0] \vdash \mathbf{e}_1 : \tau_1}{\Gamma \vdash \{\text{rec } \{\langle \text{id} \rangle : \tau_0 \ \mathbf{e}_0\} \ \mathbf{e}_1\} : \tau_1}$$

TRCFAE Type Checker

```
(define typecheck : (FAE TypeEnv -> Type)
  (lambda (fae env)
    (type-case FAE fae
      ...
      [rec (name ty rhs-expr body-expr)
        (local [(define rhs-ty (parse-type ty))
                (define new-ds (aBind name
                                      rhs-ty
                                      env))])
          (if (equal? rhs-ty (typecheck rhs-expr new-ds))
              (typecheck body-expr new-ds)
              (type-error rhs-expr (to-string rhs-ty)))))]))
```

$$\frac{\Gamma[\langle \text{id} \rangle \leftarrow \tau_0] \vdash \mathbf{e}_0 : \tau_0 \quad \Gamma[\langle \text{id} \rangle \leftarrow \tau_0] \vdash \mathbf{e}_1 : \tau_1}{\Gamma \vdash \{\text{rec } \{ \langle \text{id} \rangle : \tau_0 \ \mathbf{e}_0 \} \ \mathbf{e}_1 \} : \tau_1}$$

Variants

- A *grade* is either
 - *false*
 - *num*
- A *move* is either
 - *(left num)*
 - *(right num)*
- An *expr* is either
 - *(num num)*
 - *(app expr expr)*
 - ...

Variants

```
{with {left : (num -> (num * num))
      {fun {x : num}
          {pair 0 x}}}}
{with {right : (num -> (num * num))
      {fun {x : num}
          {pair 1 x}}}}
{with {displacement : ((num * num) -> num)
      {fun {p : (num * num)}
          {if0 {fst p}
                {- 0 {snd p}}
                {snd p}}}}
      {displacement {left 5}}}}}}
```

Variants

```
{with {grade : (num -> (num * (num * bool)))
      {fun {x : num}
          {pair 0 {pair x false}}}}
  {with {pf : (num -> (num * (num * bool)))
        {fun {y : bool}
            {pair 1 {pair 0 y}}}}
    {with {pass? : ((num * (num * bool)) -> bool)
          {fun {p : (num * (num * bool))}
              {if0 {fst p}
                   {> {fst {snd p}} 70}
                   {snd {snd p}}}}}}
      {pass? {grade 96}}}}}}
```

Have to make up a value for the other type, but this can be made to work always using thunks

Recursive Datatypes

```
{with {empty : (num * ...)  
      {pair 0 ...}}  
  {with {cons : (num -> ((num * ...) -> (num * ...)))  
        {fun {x : num}  
          {fun {r : (num * ...)}  
            {pair 1 {pair x r}}}}}}  
    {{cons 1} {{cons 2} {{cons 3} empty}}}}}}
```

Stuck again with ...

Recursive Datatypes

Add `withtype` and `cases`:

```
{withtype {numlist {empty num}
            {cons (num * numlist)}}}
{rec {len : (numlist -> num)
      {fun {l : numlist}
          {cases numlist l
              {empty {n} 0}
              {cons {fxr} {+ 1 {len {snd fxr}}}}}}}}
{len {cons {pair 1 {cons {pair 2 {empty 0}}}}}}}}
```


TVRCFAE Grammar

```
<TVRCFAE> ::= <num>
| {+ <TVRCFAE> <TVRCFAE>}
| {- <TVRCFAE> <TVRCFAE>}
| <id>
| {fun {<id>} <TVRCFAE>}
| {<TVRCFAE> <TVRCFAE>}
| {if0 <TVRCFAE> <TVRCFAE> <TVRCFAE>}
| {rec {<id> : <TE> <TVRCFAE>} <TVRCFAE>}
| {withtype {<tyid> {<id> <TE>}
              {<id> <TE>}}
      <TVRCFAE>}
| {cases <tyid> <TVRCFAE>
      {<id> {<id>} <TVRCFAE>}
      {<id> {<id>} <TVRCFAE>}}
```

NEW

NEW

```
<TE> ::= num
| (<TE> -> <TE>)
| <tyid>
```

NEW

Well-Formed Type Expressions

- Might be ok:

```
{withtype {fruit {apple num}
           {banana (num -> num) }}
 ... {fun {x : fruit} ...} ...}
```

- Not ok:

```
{fun {x : fruit} ...}
```

$$\Gamma \vdash \text{num} \qquad \frac{\Gamma \vdash \tau_1 \quad \Gamma \vdash \tau_2}{\Gamma \vdash (\tau_1 \rightarrow \tau_2)}$$

[... $\langle \text{tyid} \rangle = \langle \text{id} \rangle_1 @ \tau_1 + \langle \text{id} \rangle_2 @ \tau_2$...] $\vdash \langle \text{tyid} \rangle$

TVRCFAE Type Checker

$$\Gamma' = \Gamma[\langle \text{tyid} \rangle = \langle \text{id} \rangle_1 @ \tau_1 + \langle \text{id} \rangle_2 @ \tau_2, \langle \text{id} \rangle_1 \leftarrow (\tau_1 \rightarrow \langle \text{tyid} \rangle), \langle \text{id} \rangle_2 \leftarrow (\tau_2 \rightarrow \langle \text{tyid} \rangle)]$$

$$\Gamma' \vdash \tau_1 \quad \Gamma' \vdash \tau_2 \quad \Gamma' \vdash e : \tau_0$$

$$\Gamma \vdash \{ \text{withtype } \{ \langle \text{tyid} \rangle \{ \langle \text{id} \rangle_1 \ \tau_1 \} \{ \langle \text{id} \rangle_2 \ \tau_2 \} \} e \} : \tau_0$$

$$\Gamma' = \Gamma[\langle \text{tyid} \rangle = \langle \text{id} \rangle_1 @ \tau_1 + \langle \text{id} \rangle_2 @ \tau_2]$$

$$\Gamma' \vdash e_0 : \langle \text{tyid} \rangle \quad \Gamma' [\langle \text{id} \rangle_3 \leftarrow \tau_1] \vdash e_1 : \tau_0 \quad \Gamma' [\langle \text{id} \rangle_4 \leftarrow \tau_2] \vdash e_2 : \tau_0$$

$$\Gamma' \vdash \{ \text{cases } \langle \text{tyid} \rangle e_0 \{ \langle \text{id} \rangle_1 \{ \langle \text{id} \rangle_3 \} e_1 \} \{ \langle \text{id} \rangle_2 \{ \langle \text{id} \rangle_4 \} e_2 \} \} : \tau_0$$

Warning: later, we'll discuss why the `withtype` rule is not quite right

TVRCFAE Expression Datatypes

```
(define-type FAE
  [with-type (name : symbol)
             (var1-name : symbol)
             (var1-ty : TE)
             (var2-name : symbol)
             (var2-ty : TE)
             (body-expr : FAE)]
  [cases (name : symbol)
         (dispatch-expr : FAE)
         (var1-name : symbol)
         (bind1-name : symbol)
         (rhs1-expr : FAE)
         (var2-name : symbol)
         (bind2-name : symbol)
         (rhs2-expr : FAE)])

(define-type TE
  ...
  [idTE (name : symbol)])
```

TVRCFAE Value and Environment Datatypes

```
(define-type FAE-Value
  ...
  [variantV (right? : boolean)
            (val : FAE-Value)]
  [constructorV (right? : boolean)])
```

```
(define-type TypeEnv
  ...
  [tBind (name : symbol)
        (var1-name : symbol)
        (var1-type : Type)
        (var2-name : symbol)
        (var2-type : Type)
        (rest : TypeEnv)])
```

TVRCFAE Interpreter

```
(define (interp a-fae ds)
  (type-case FAE a-fae
    ...
    [with-type (type-name var1-name var1-te
                    var2-name var2-te
                    body-expr)
      (interp body-expr
              (aSub var1-name
                    (constructorV false)
                    (aSub var2-name
                          (constructorV true)
                          ds))))])
  ...))
```

TVRCFAE Interpreter

```
(define (interp a-fae ds)
  (type-case FAE a-fae
    ...
    [app (fun-expr arg-expr)
      (local [(define fun-val
                 (interp fun-expr ds))
              (define arg-val
                 (interp arg-expr ds))]
        (type-case FAE-Value fun-val
          [closureV (param body ds)
            (interp body
                      (aSub param
                           arg-val
                           ds))]
          [constructorV (right?)
            (variantV right? arg-val)]
          [else (error 'interp "not applicable")]]))]
    ...))
```

TVRCFAE Interpreter

```
(define (interp a-fae ds)
  (type-case FAE a-fae
    ...
    [cases (ty dispatch-expr
             var1-name var1-id var1-rhs
             var2-name var2-id var2-rhs)
           (type-case FAE-Value (interp dispatch-expr ds)
             [variantV (right? val)
                      (if (not right?)
                          (interp var1-rhs (aSub var1-id
                                                  val
                                                  ds))
                          (interp var2-rhs (aSub var2-id
                                                  val
                                                  ds)))]
             [else (error 'interp "not a variant result")])]
    ...))
```


TVRCFAE Type Lookup

```
(define (get-type name-to-find env)
  (type-case TypeEnv env
    [mtEnv () (error 'get-type "free variable, so no type")]
    [aBind (name ty rest)
      (if (symbol=? name-to-find name)
          ty
          (get-type name-to-find rest))]
    [tBind (name var1-name var1-ty var2-name var2-ty rest)
      (get-type name-to-find rest)]))
```

TVRCFAE Type Lookup

```
(define (find-type-id name-to-find env)
  (type-case TypeEnv env
    [mtEnv () (error 'get-type "free type name, so no type")]
    [aBind (name ty rest)
      (find-type-id name-to-find rest)]
    [tBind (name var1-name var1-ty var2-name var2-ty rest)
      (if (symbol=? name-to-find name)
          env
          (find-type-id name-to-find rest))]))
```

TVRCFAE Type-Expression Checking

```
(define (validtype ty env)
  (type-case Type ty
    [numT () (mtEnv)]
    [boolT () (mtEnv)]
    [arrowT (a b) (begin
                     (validtype a env)
                     (validtype b env))]
    [idT (id) (find-type-id id env)]))
```

TVRCFAE Type Checking

```
(define typecheck : (FAE TypeEnv -> Type)
  (lambda (fae env)
    (type-case FAE fae
      ...
      [fun (name te body)
        (local [(define arg-type (parse-type te))]
          (begin
            (validtype arg-type env)
            (arrowT arg-type
                    (typecheck body (aBind name
                                           arg-type
                                           env))))))])
    ...)))
```

TVRCFAE Type Checking

```
(define typecheck : (FAE TypeEnv -> Type)
  (lambda (fae env)
    (type-case FAE fae
      ...
      [with-type (type-name var1-name var1-te var2-name var2-te
                    body-expr)
        (local [(define var1-ty (parse-type var1-te))
                (define var2-ty (parse-type var2-te))
                (define new-env (tBind type-name
                                       var1-name var1-ty
                                       var2-name var2-ty
                                       env))])
          (begin
            (validtype var1-ty new-env)
            (validtype var2-ty new-env)
            (typecheck body-expr
                       (aBind var1-name
                              (arrowT var1-ty
                                       (idT type-name))
                              (aBind var2-name
                                     (arrowT var2-ty
                                             (idT type-name))
                                     new-env))))))
    ...)))
```

TVRCFAE Type Checking

```
(define typecheck : (FAE TypeEnv -> Type)
  (lambda (fae env)
    (type-case FAE fae
      ...
      [cases (type-name dispatch-expr
                var1-name var1-id var1-rhs
                var2-name var2-id var2-rhs)
        (local [(define bind (find-type-id type-name env))]
          (if (and (equal? var1-name (tBind-var1-name bind))
                  (equal? var2-name (tBind-var2-name bind)))
              (type-case Type (typecheck dispatch-expr env)
                [idT (name)
                  (if (equal? name type-name)
                      (local [(define rhs1-ty
                                (typecheck var1-rhs
                                           (aBind var1-id
                                                (tBind-var1-type bind)
                                                env)))
                              (define rhs2-ty
                                (typecheck var2-rhs
                                           (aBind var2-id
                                                (tBind-var2-type bind)
                                                env)))]
                        (if (equal? rhs1-ty rhs2-ty)
                            rhs1-ty
                            (type-error var2-rhs (to-string rhs1-ty)))
                        (type-error dispatch-expr (to-string type-name))))
                  [else (type-error dispatch-expr (to-string type-name))])
              (type-error fae "matching variant names")))]
      ...)))
```

Type Soundness

Type soundness is a theorem of the form

If $\emptyset \vdash e : \tau$, then running e never produces an error

If we add division, then divide-by-zero errors may be ok:

If $\emptyset \vdash e : \tau$, then running e never produces an error except divide-by-zero

In general, soundness rules out a certain class of run-time errors

Soundness fails \Rightarrow bug in type rules

Type Soundness in TRCFAE

TRCFAE has a bug:

```
{rec {f : (num -> num)
      f}
 {f 10}}
```

One solution: adjust the soundness theorem to allow a run-time error

Another solution: change the grammar for **rec**

```
<TIFAE> ::= ...
          | {rec {<id> : <tyexp>
                  {fun {<id> : <tyexp>
                        <TIFAE>}}}
            <TIFAE>}
```


Type Soundness in TVRCFAE

TCRCFAE has a bug, too:

```

{{withtype {foo {a num} {b num}}
  {fun {x : foo} {+ {cases foo x
                    {a {n} n}
                    {b {n} n}}}}}
{withtype {foo {c (num -> num)} {d num}}
  {c {fun {y : num} y}}}}

```

Solution 1: no local type declarations

Solution 2: don't let `<tyid>` escape `withtype`

$$\Gamma' = \Gamma[\langle \text{tyid} \rangle = \langle \text{id} \rangle_1 @ \tau_1 + \langle \text{id} \rangle_2 @ \tau_2, \langle \text{id} \rangle_1 \leftarrow (\tau_1 \rightarrow \langle \text{tyid} \rangle), \langle \text{id} \rangle_2 \leftarrow (\tau_2 \rightarrow \langle \text{tyid} \rangle)]$$

`<tyid>` not in τ_0

$$\Gamma' \vdash \tau_1 \quad \Gamma' \vdash \tau_2 \quad \Gamma' \vdash e : \tau_0$$

$$\Gamma \vdash \{\text{withtype} \{ \langle \text{tyid} \rangle \{ \langle \text{id} \rangle_1 \ \tau_1 \} \{ \langle \text{id} \rangle_2 \ \tau_2 \} \} e \} : \tau_0$$