

# Part I

# Allocation

Constructor calls are allocation:

```
(define (interp)
  (type-case ExprD expr-reg
    ...
    [lamD (body-expr)
      (begin
        (set! v-reg (closV body-expr env-reg))
        (continue))]
    ...))
```

```
(define (continue)
  ...
  [addSecondK (r env k)
    (begin
      (set! expr-reg r)
      (set! env-reg sc)
      (set! k-reg (doAddK v-reg k))
      (interp))]
  ...)
```

# Deallocation

Where does **free** go?

```
(define (continue)
  ...
  [doAddK (v1 k)
    (begin
      (set! v-reg (num+ v1 v-reg))
      (free k-reg) ; ???
      (set! k-reg k)
      (continue))]
  ...
  [doAppK (fun-val k)
    (begin
      (set! expr-reg (closV-body fun-val))
      (set! env-reg (cons v-reg
                          (closV-env fun-val)))
      (set! k-reg k)
      (free fun-val) ; ???
      (interp))]
  ...)
```

# Deallocation

```
[doAddK (v1 k)
  (begin
    (set! v-reg (num+ v1 v-reg))
    (free k-reg) ; ???
    (set! k-reg k)
    (continue)))]
```

- Without `let/cc`, this `free` is fine, because the continuation can't be referenced anywhere else
- A continuation object is always freed as `(free k-reg)`, which is why many language implementations use a stack

# Deallocation

```
[doAppK (fun-val k)
  (begin
    (set! expr-reg (closV-body fun-val))
    (set! env-reg (cons v-reg
                        (closV-env fun-val)))
    (set! k-reg k)
    (free fun-val) ; ???
    (interp))] ]
```

- This free is *not* ok, because the closure might be kept in a environment somewhere
- Need to free only if no one else is using it...

# Code and Data

An **object** is any record allocated during **interp** and **continue**

Assume that expressions are allocated “statically”

- **compile** uses **code-malloc1**, etc.
- Only try to free objects allocated during **interp** and **continue**

## Part 2

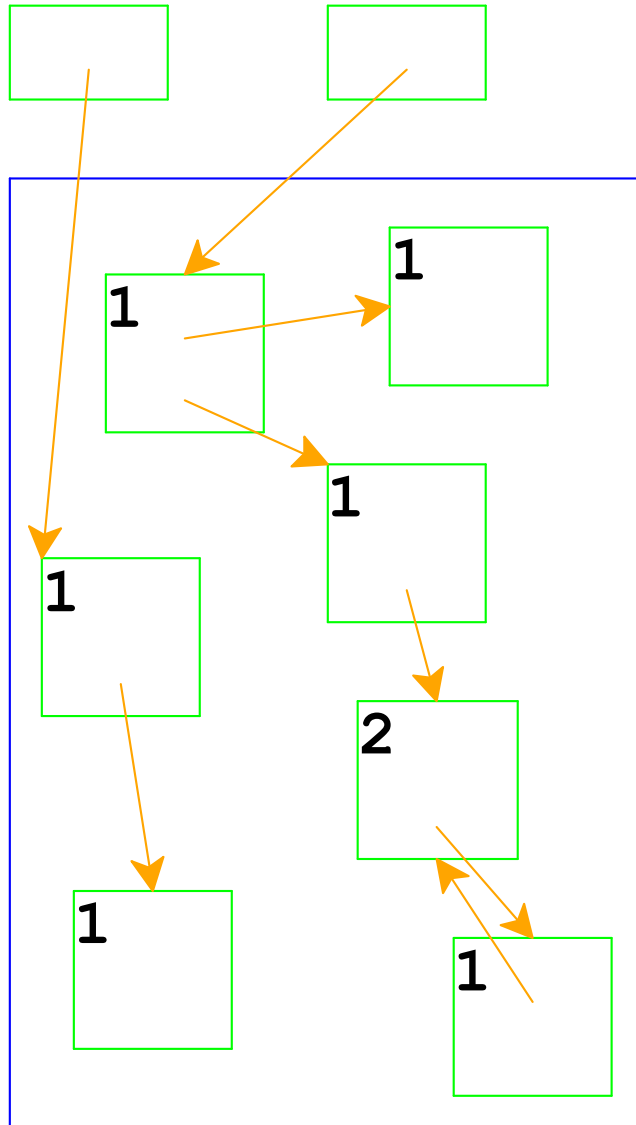
# Reference Counting

**Reference counting:** a way to know whether an object has other users

- Attach a count to every object, starting at 0
- When installing a pointer to an object (into a register or another object), increment its count
- When replacing a pointer to an object, decrement its count
- When a count is decremented to 0, decrement counts for other objects referenced by the object, then free



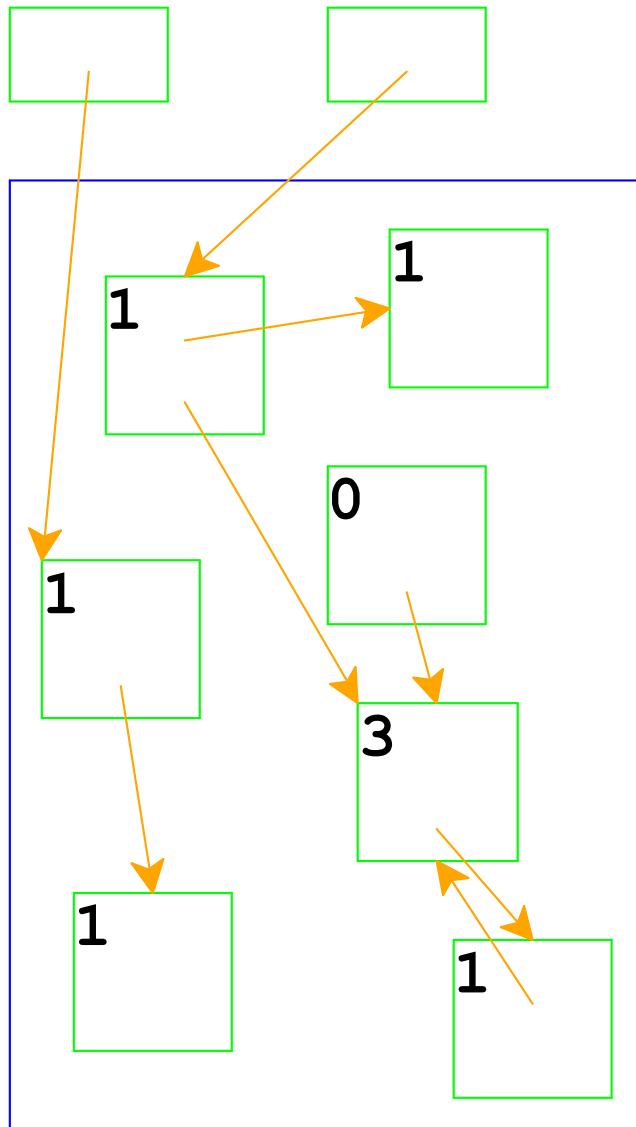
# Reference Counting



Top boxes are the registers  
**k-reg**, **v-reg**, etc.

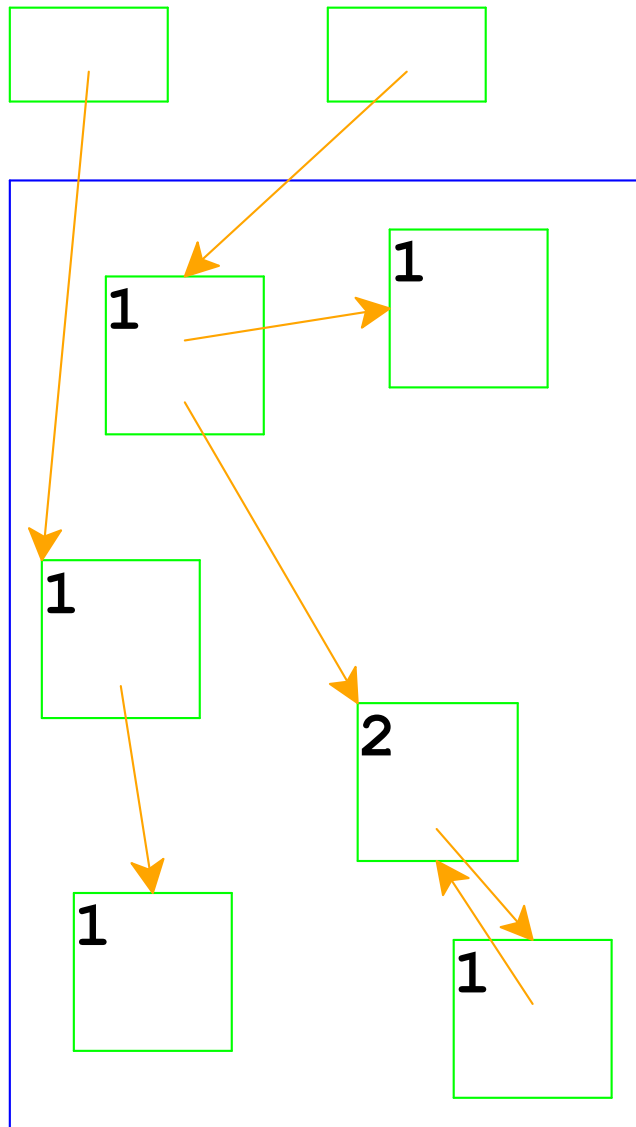
Boxes in the blue area are  
allocated with **malloc**

# Reference Counting



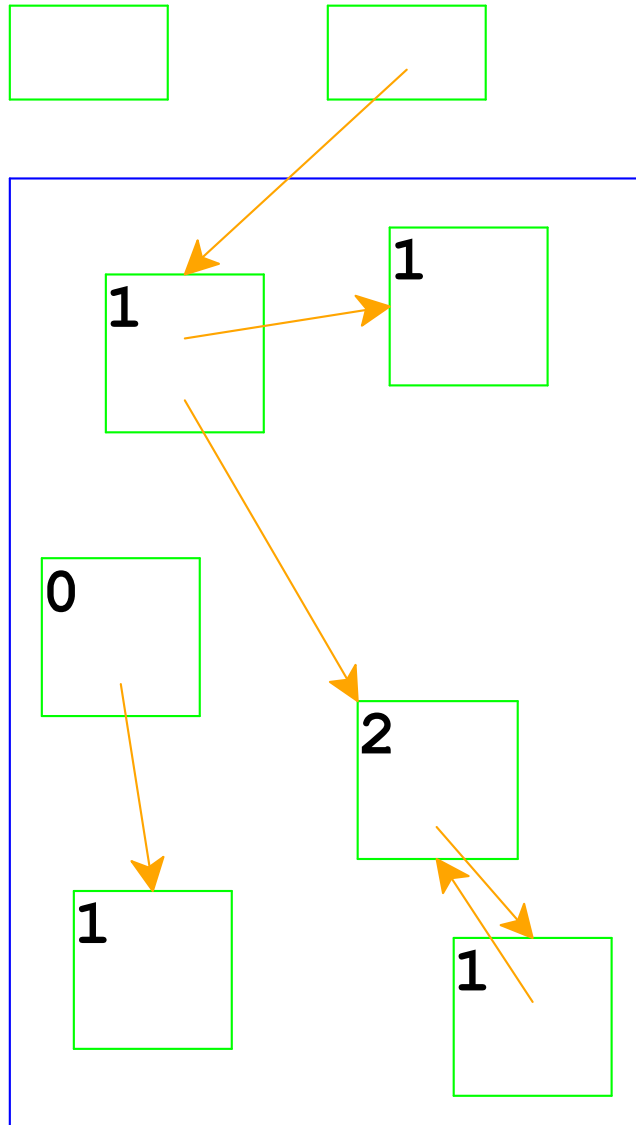
Adjust counts when a pointer is changed...

# Reference Counting



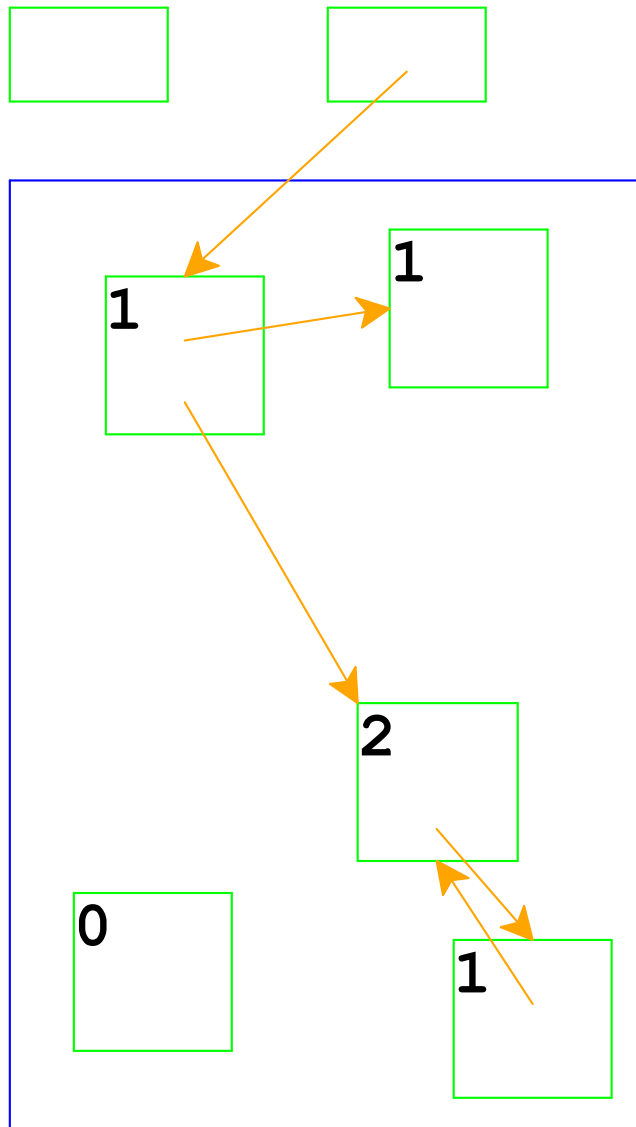
... freeing an object if its count goes to 0

# Reference Counting



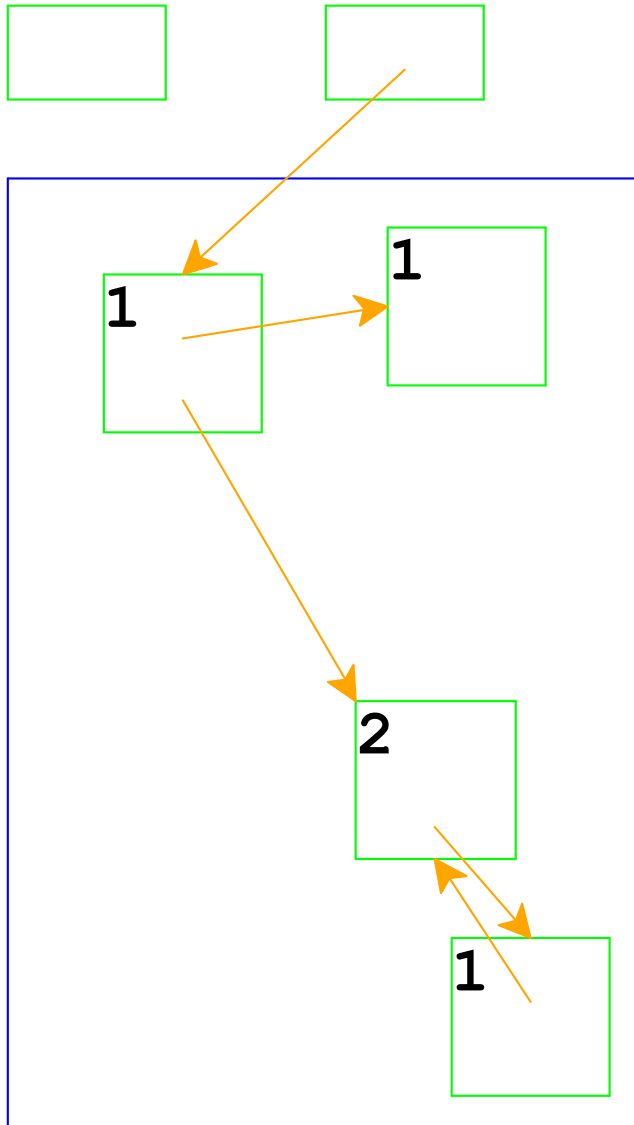
Same if the pointer is in a register

# Reference Counting



Adjust counts after frees, too...

# Reference Counting



... which can trigger more frees

# Reference Counting in an Interpreter

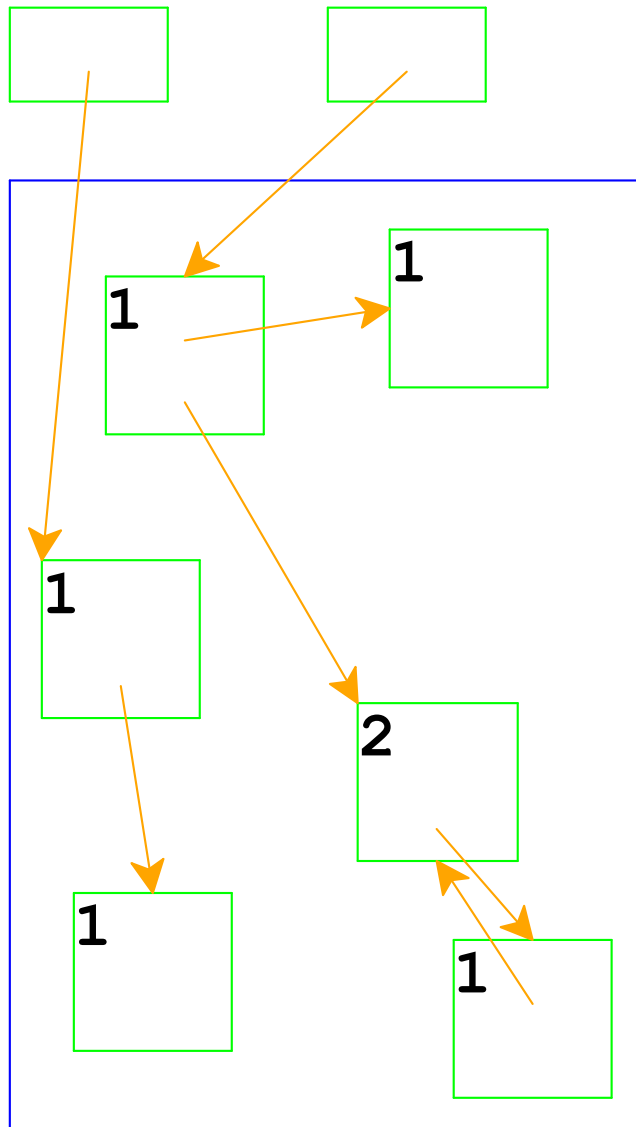
...

```
[lamC (body-expr)
  (begin
    (ref- v-reg)
    (set! v-reg
      ; must ref+ env:
      (closV body-expr env-reg))
    (ref+ v-reg)
    (continue))]
```

...

```
[doAppK (fun-val k)
  (begin
    (set! fae-reg (closV-body fun-val)) ; code is static
    (ref- env-reg)
    (set! env-reg
      ; must ref+ each arg:
      (cons v-reg (closV-env fun-val)))
    (ref+ env-reg) ; => ref+ on v-reg
    (ref+ k)
    (ref- k-reg) ; => ref- on fun-val and k
    (set! k-reg k)
    (interp))]
```

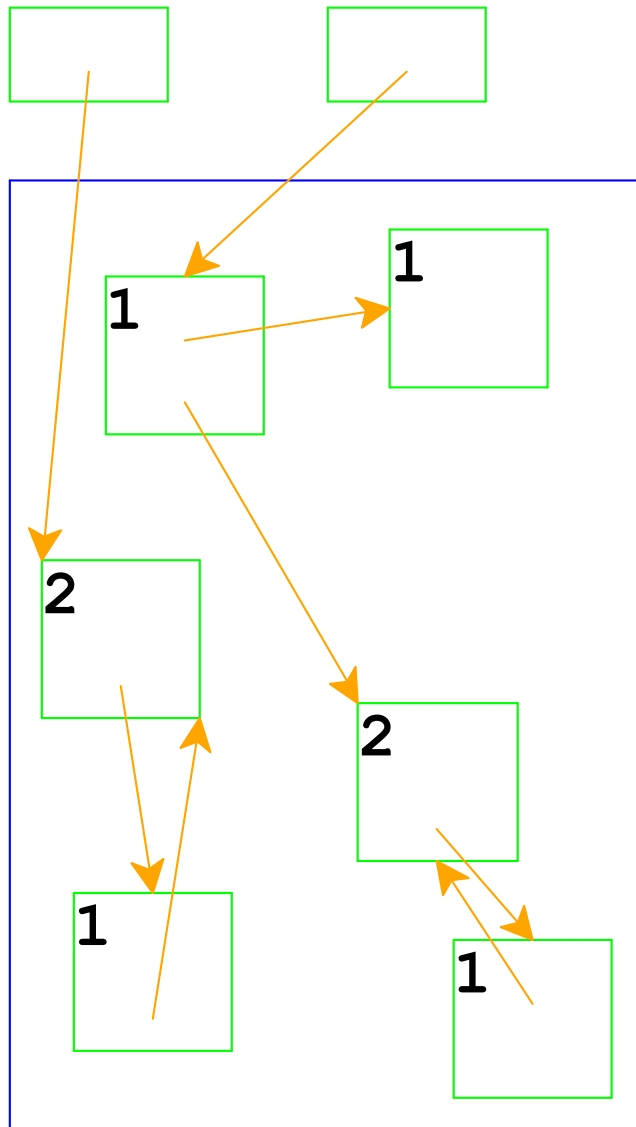
# Reference Counting And Cycles



An assignment can create a cycle...

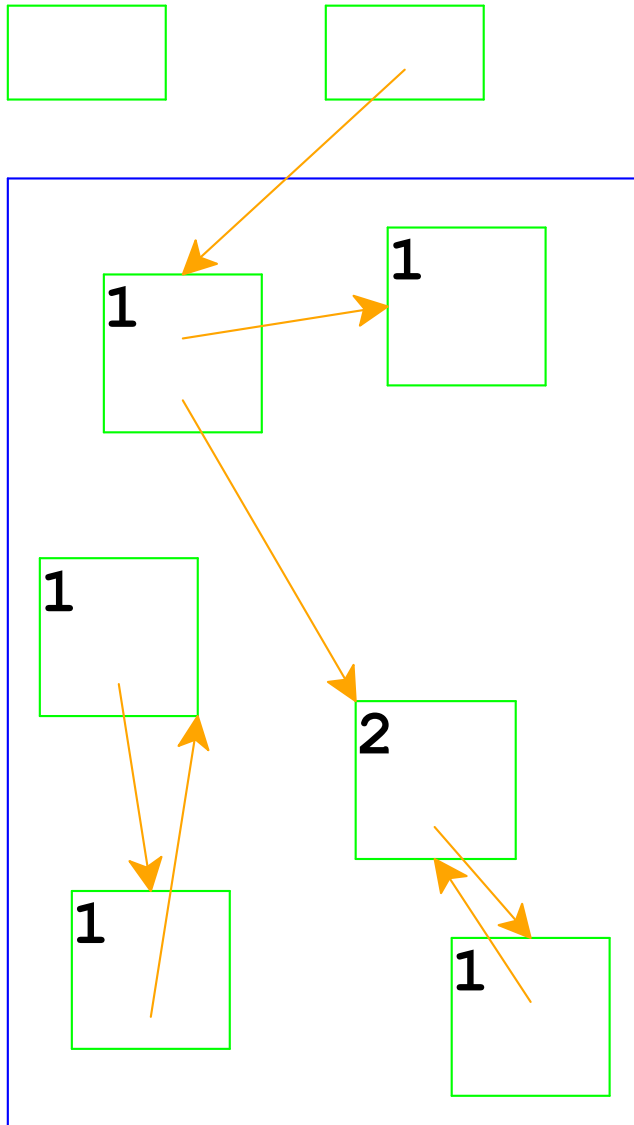


# Reference Counting And Cycles



Adding a reference increments a count

# Reference Counting And Cycles



Lower-left objects are inaccessible, but not deallocated

In general, cycles break reference counting

# Part 3

# Garbage Collection

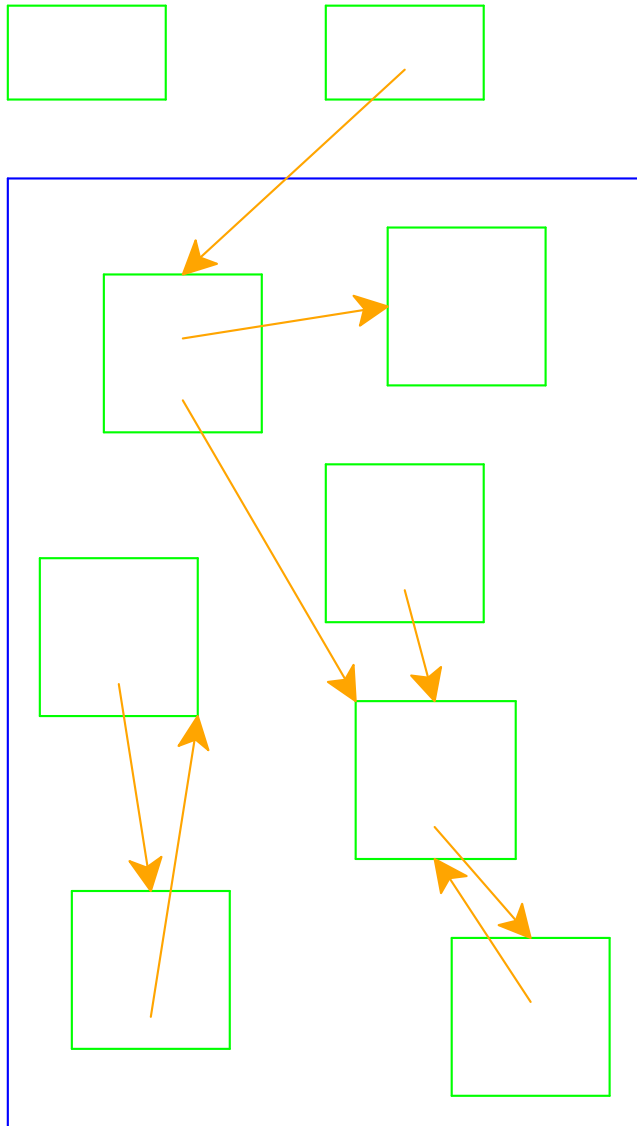
**Garbage collection:** a way to know whether an object is *accessible*

- An object referenced by a register is **live**
- An object referenced by a live object is also live
- A program can only possibly use live objects, because there is no way to get to other objects
- A garbage collector frees all objects that are not live
- Allocate until we run out of memory, then run a garbage collector to get more space

# Garbage Collection Algorithm

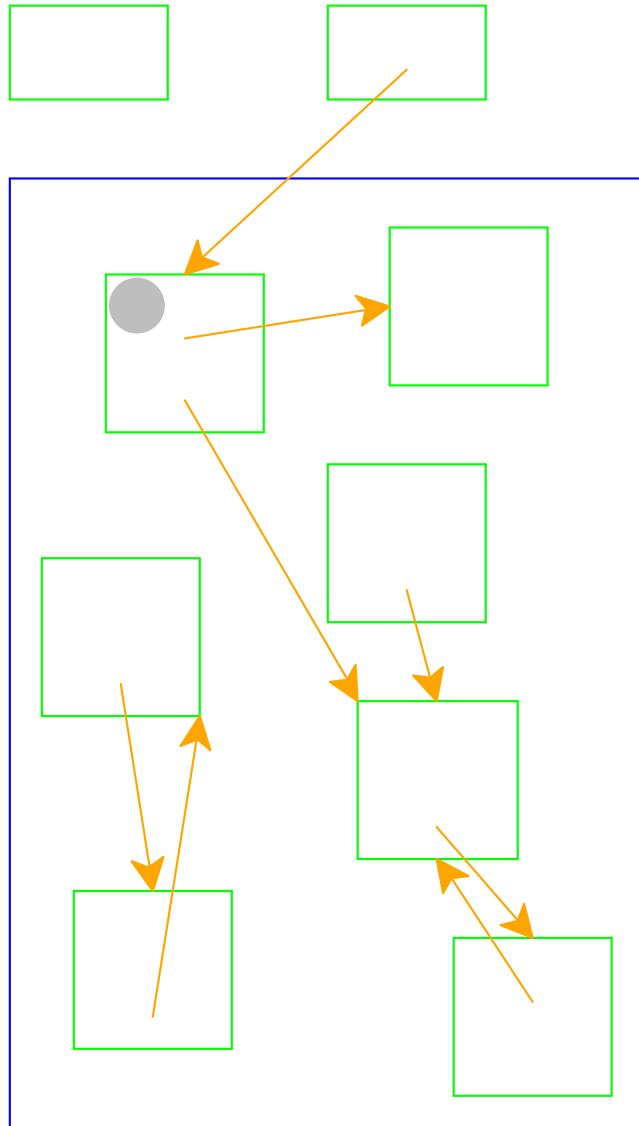
- Color all objects **white**
- Color objects referenced by registers **gray**
- Repeat until there are no gray objects:
  - Pick a gray object,  $r$
  - For each white object that  $r$  points to, make it gray
  - Color  $r$  **black**
- Deallocate all white objects

# Garbage Collection



All objects are marked white

# Garbage Collection

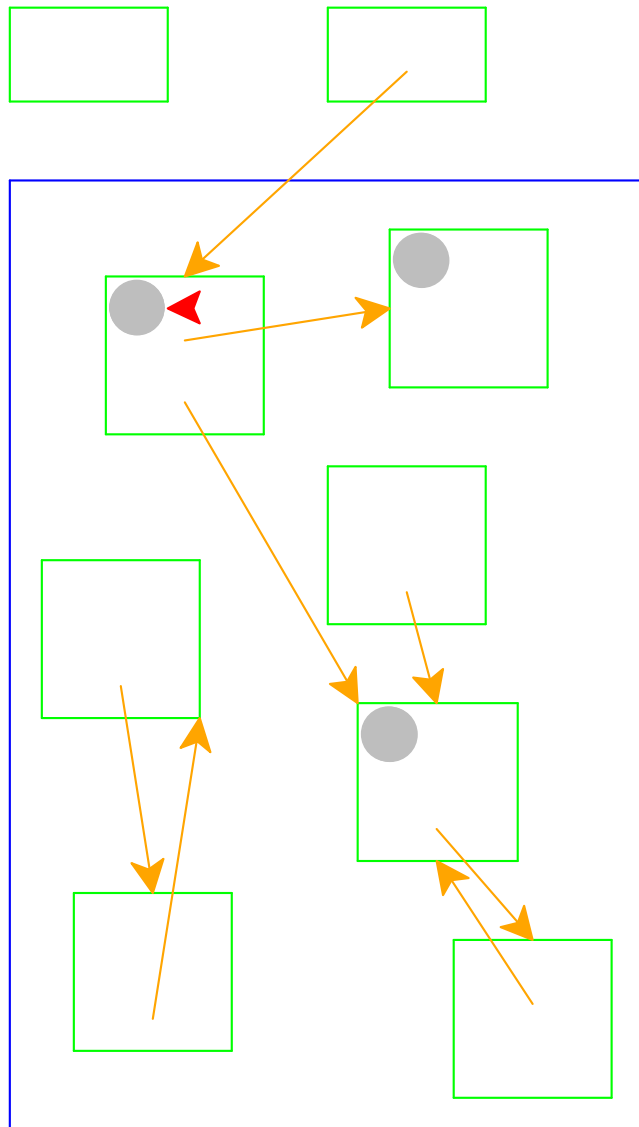


Mark objects referenced by registers as gray



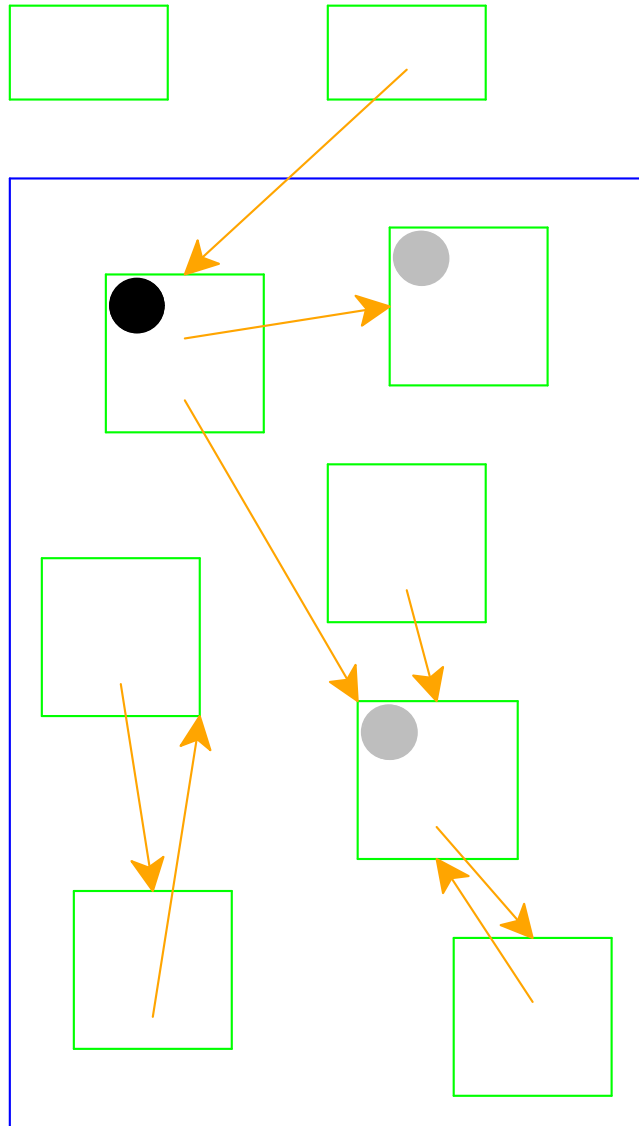


# Garbage Collection



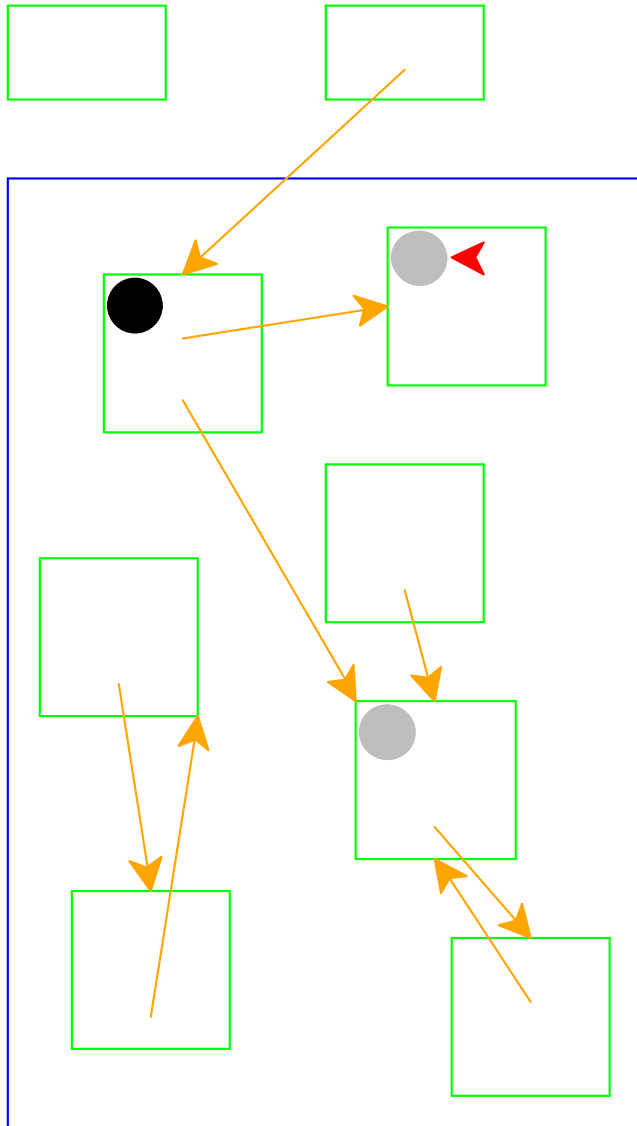
Mark white objects referenced by  
chosen object as gray

# Garbage Collection



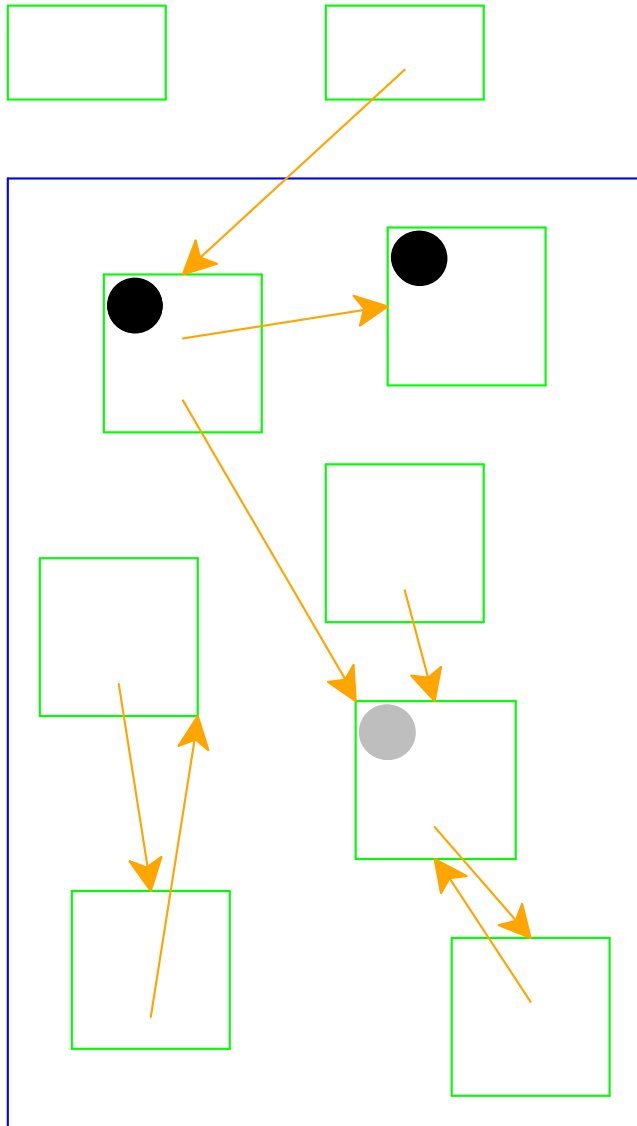
Mark chosen object black

# Garbage Collection



Start again: pick a gray object

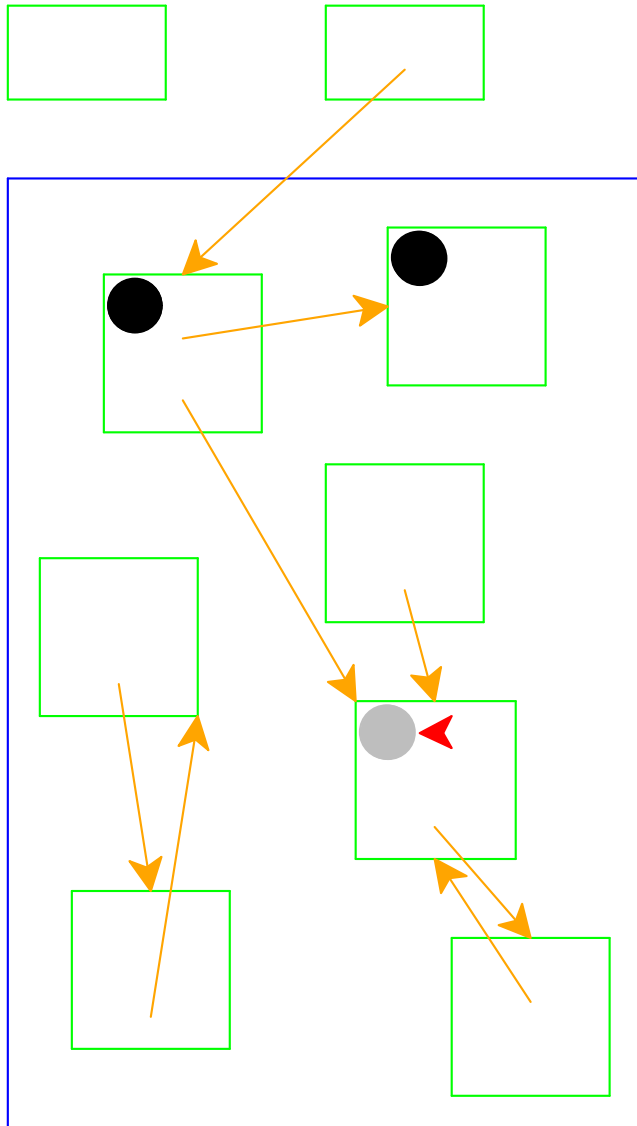
# Garbage Collection



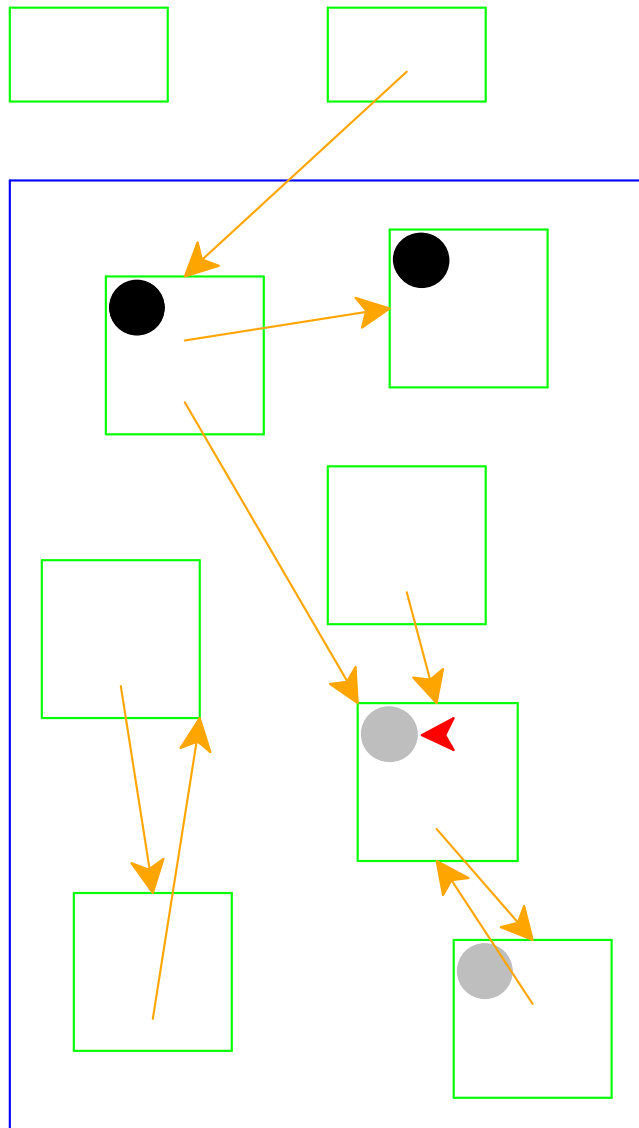
No referenced objects; mark black

# Garbage Collection

Start again: pick a gray object

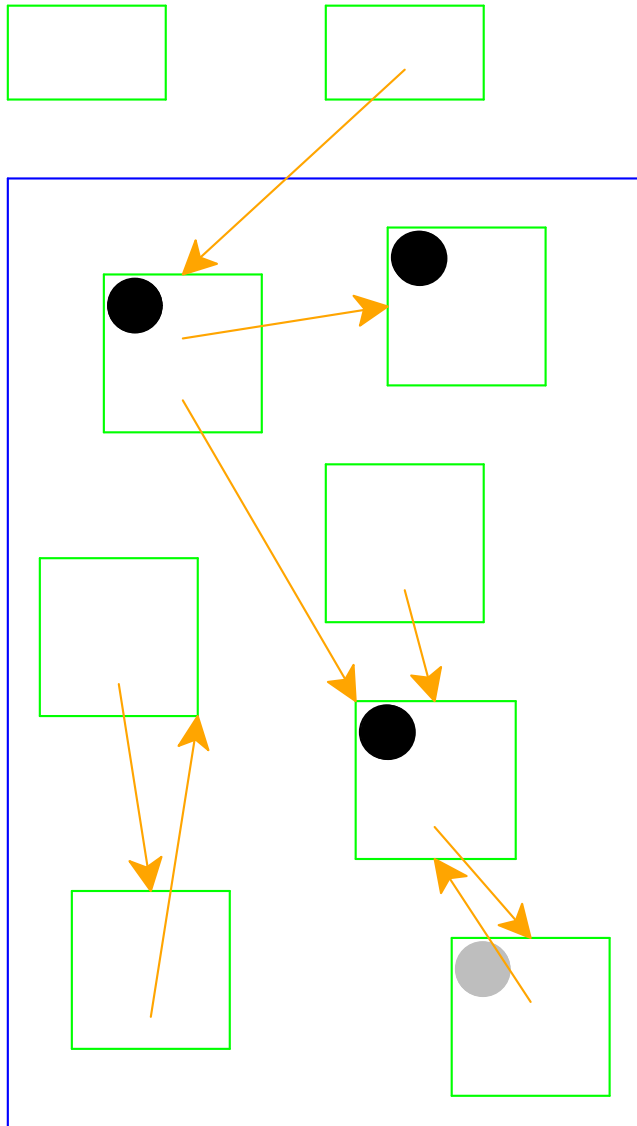


# Garbage Collection



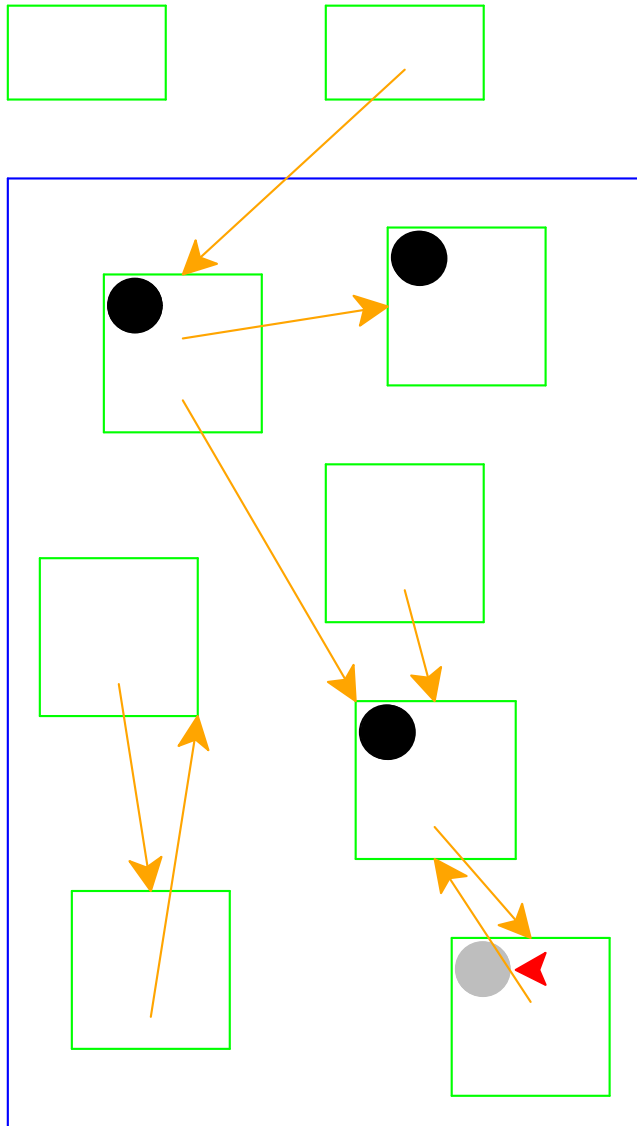
Mark white objects referenced by chosen object as gray

# Garbage Collection



Mark chosen object black

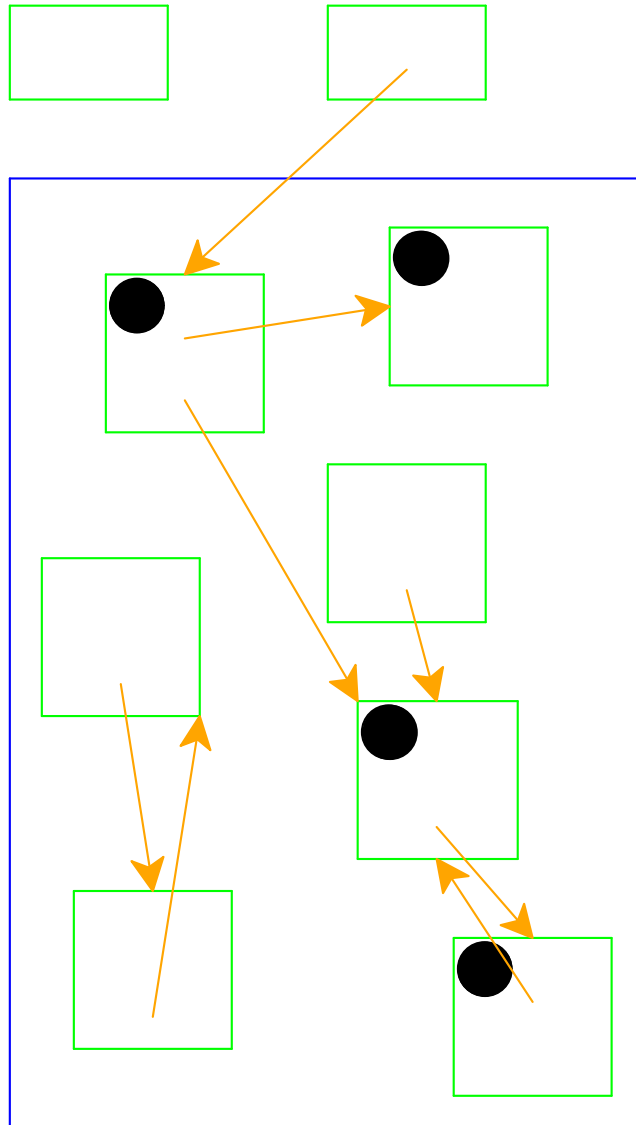
# Garbage Collection



Start again: pick a gray object

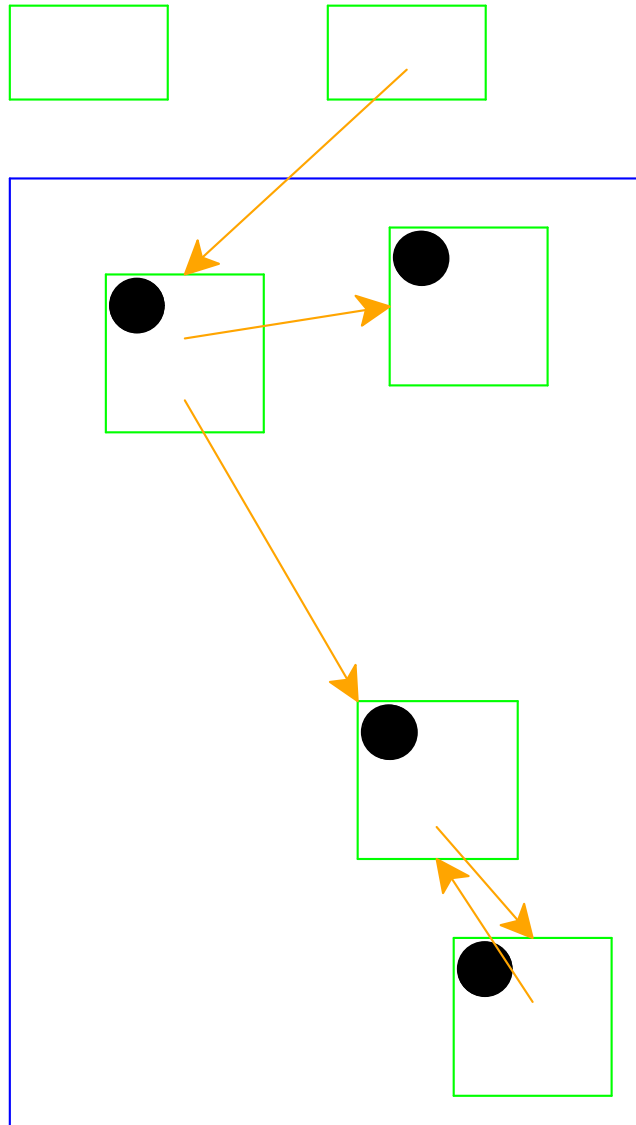


# Garbage Collection



No referenced white objects;  
mark black

# Garbage Collection



No more gray objects; deallocate white objects

Cycles **do not** break garbage collection

# Part 4

# Two-Space Copying Collectors

A **two-space** copying collector compacts memory as it collects, making allocation easier.

## Allocator:

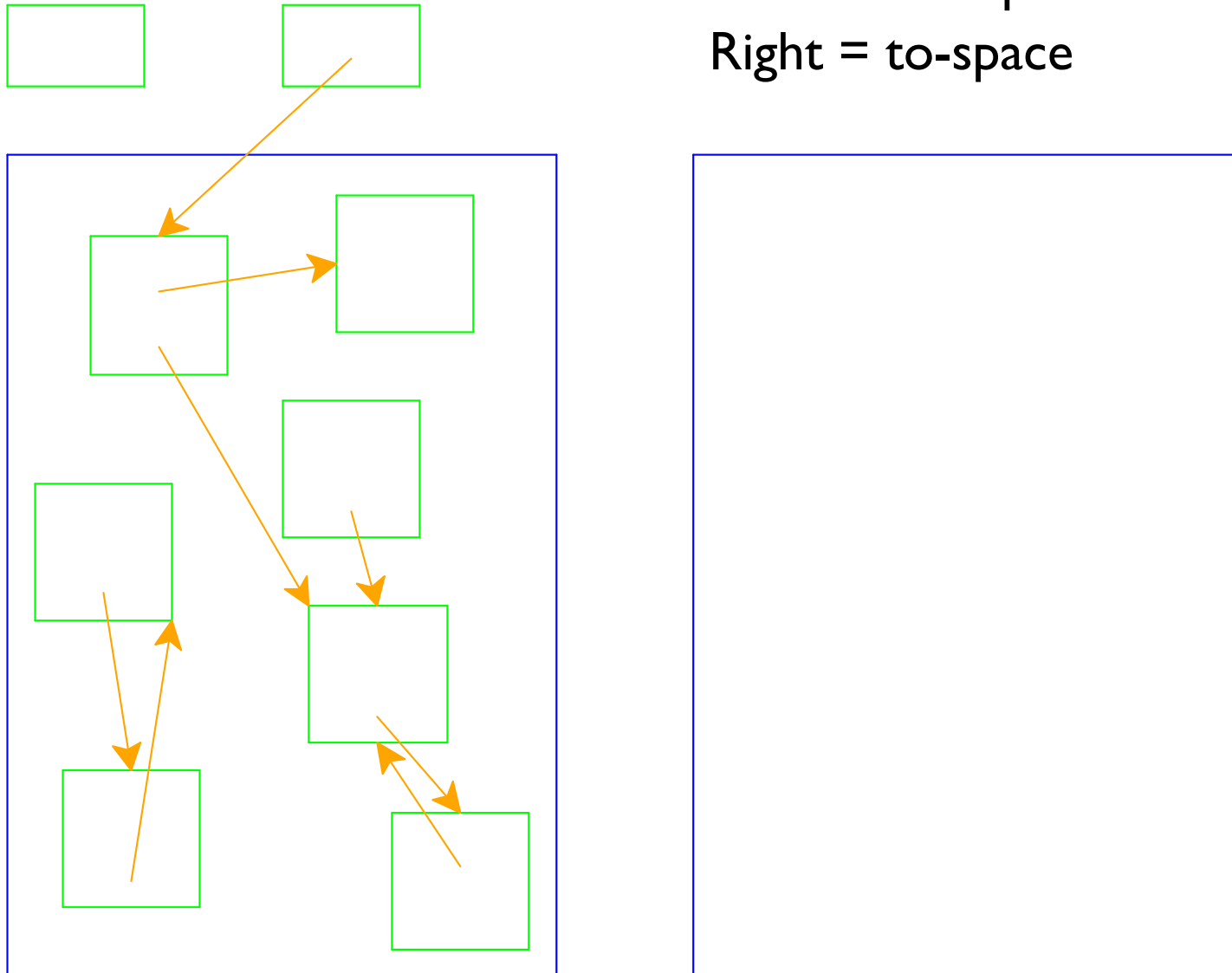
- Partitions memory into **to-space** and **from-space**
- Allocates only in **to-space**

## Collector:

- Starts by swapping **to-space** and **from-space**
- Coloring gray  $\Rightarrow$  copy from **from-space** to **to-space**
- Choosing a gray object  $\Rightarrow$  walk once through the new **to-space**, update pointers

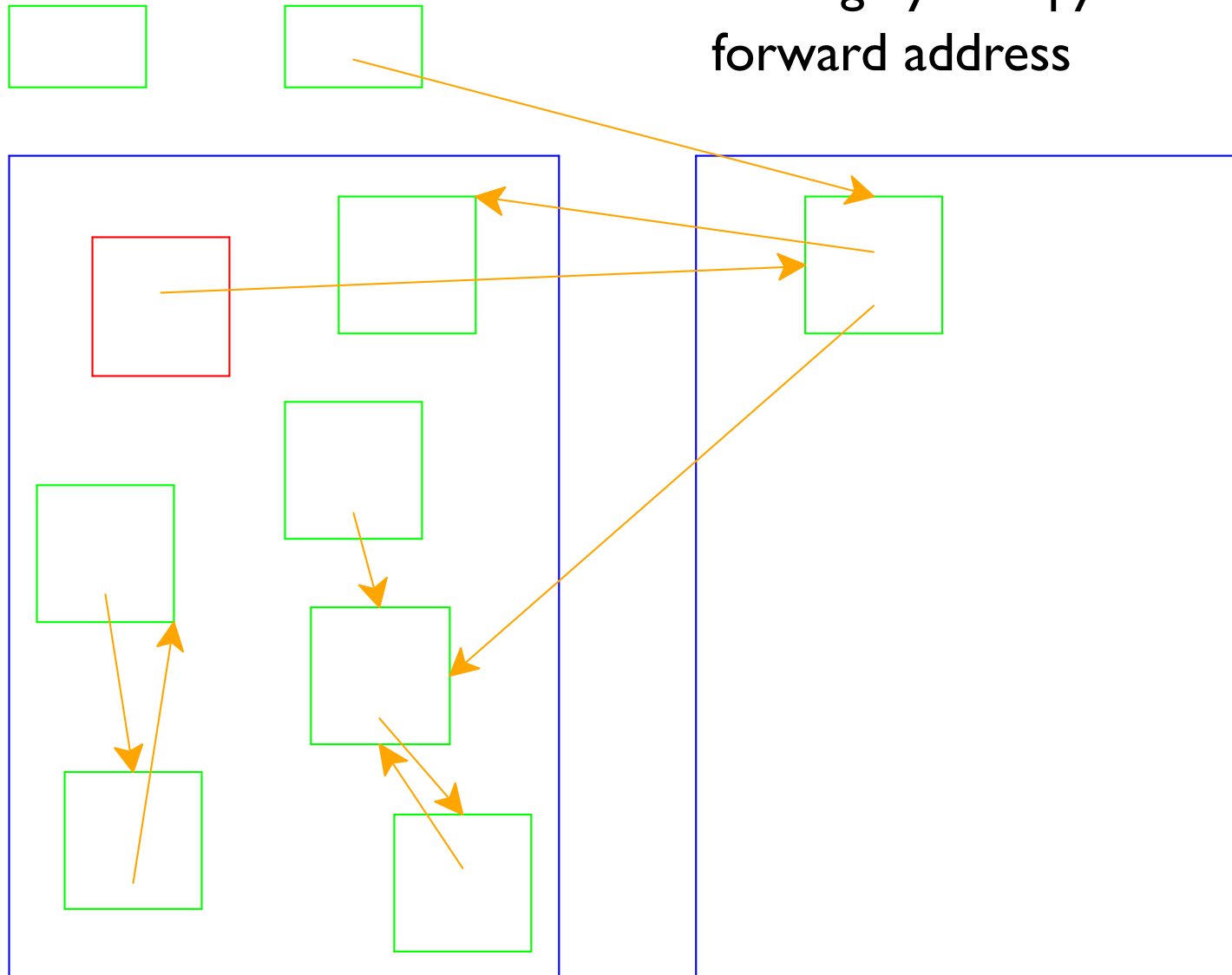
# Two-Space Collection

Left = from-space  
Right = to-space



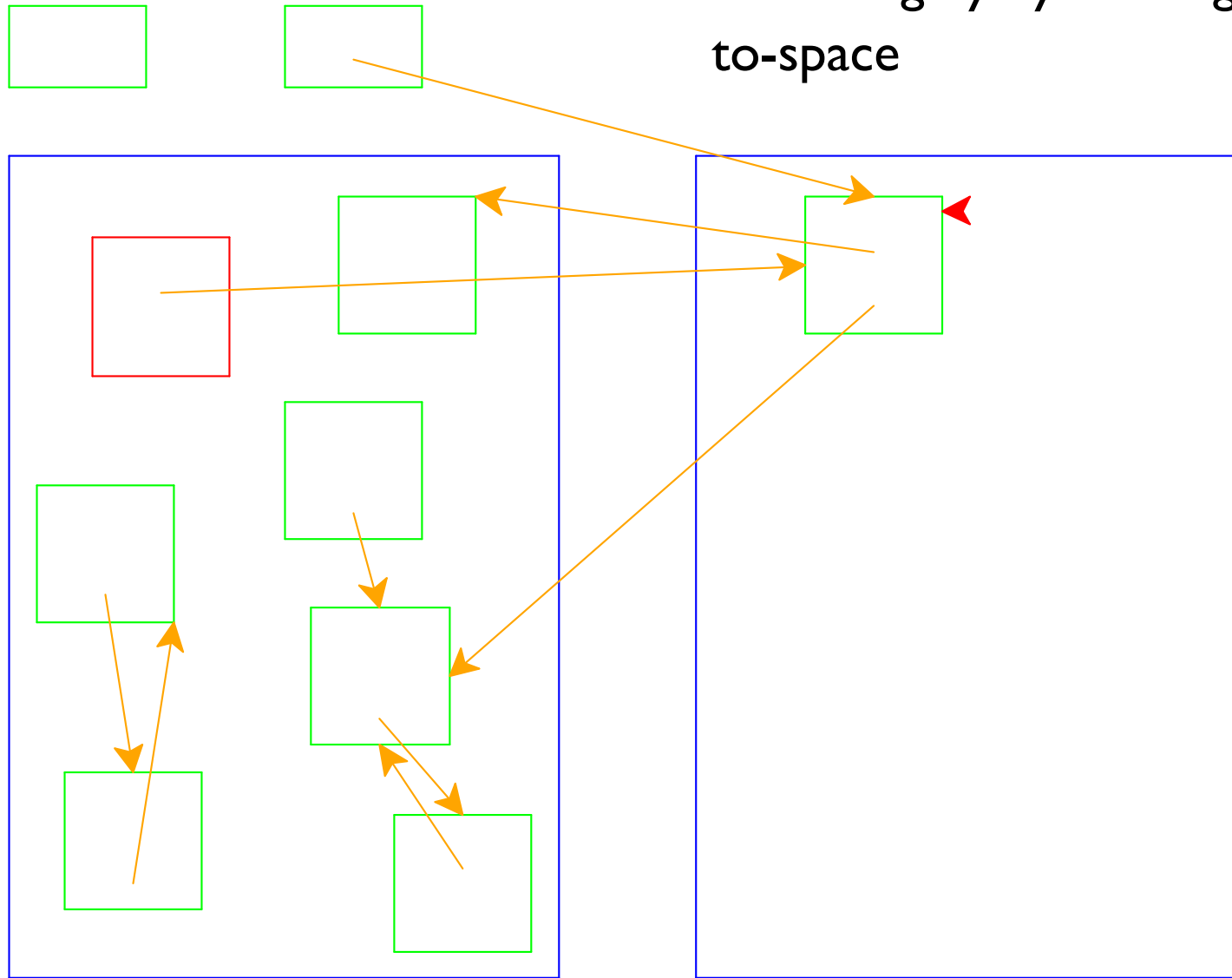
# Two-Space Collection

Mark gray = copy and leave forward address



# Two-Space Collection

Choose gray by walking through to-space

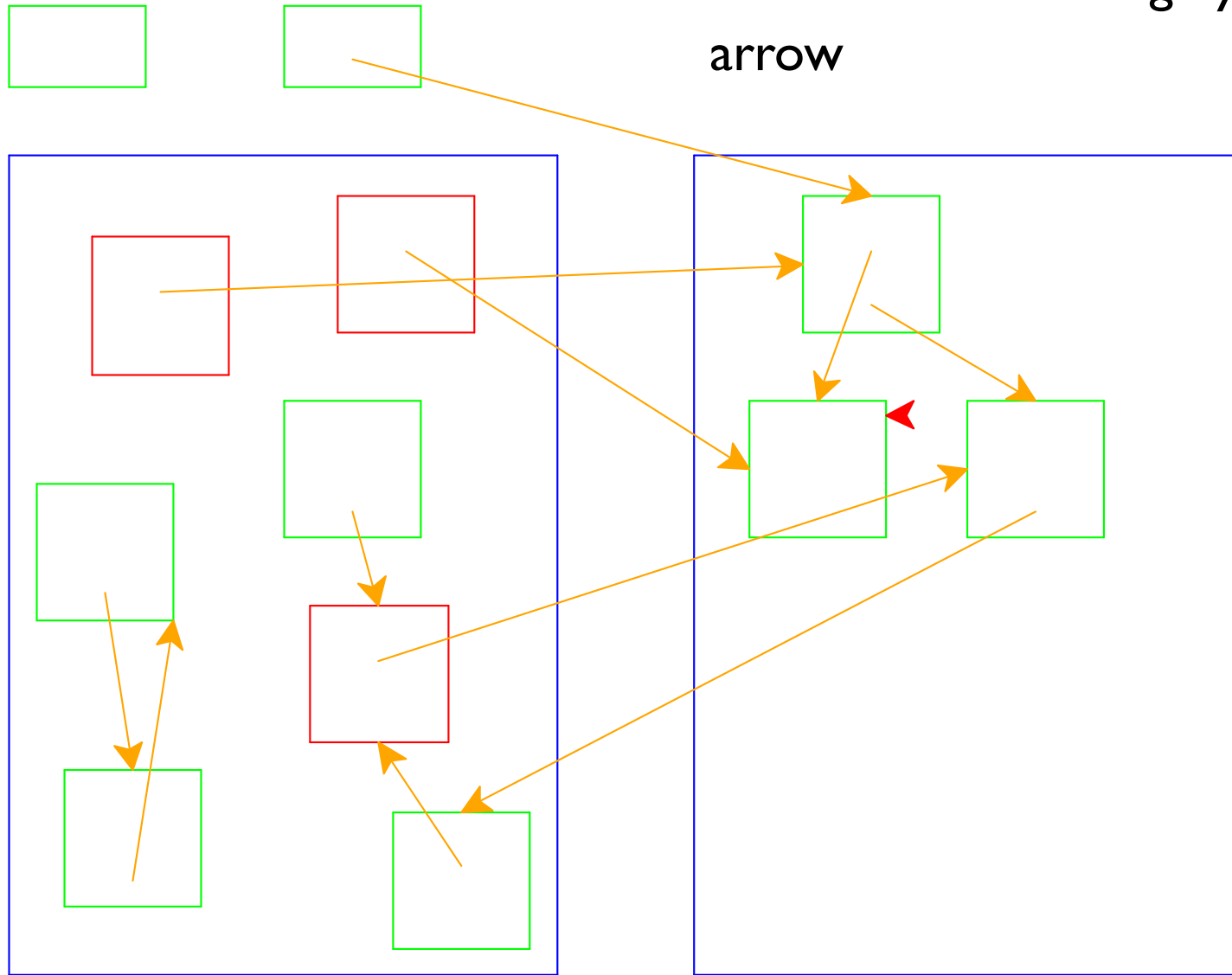






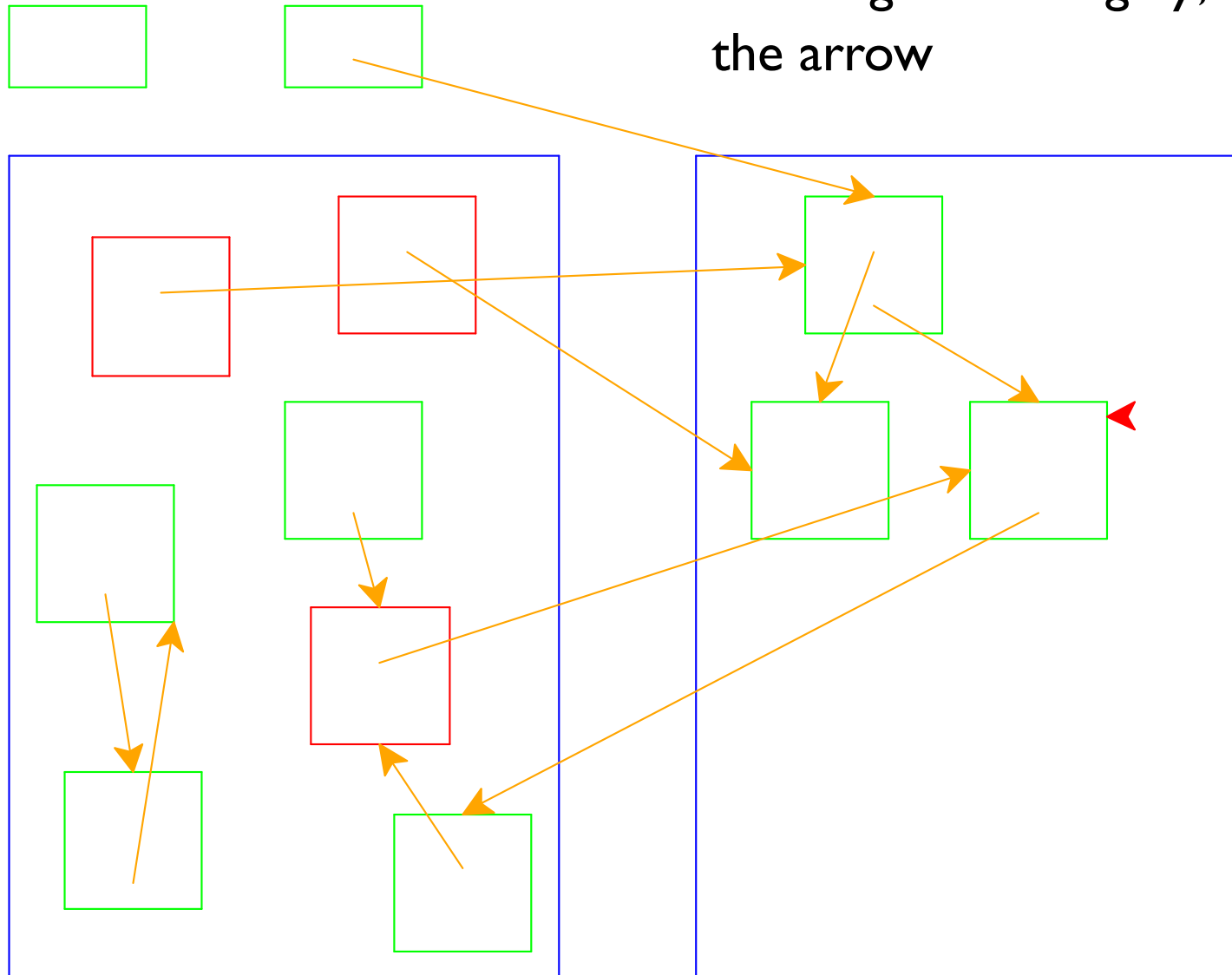
# Two-Space Collection

Mark black = move gray-choosing  
arrow



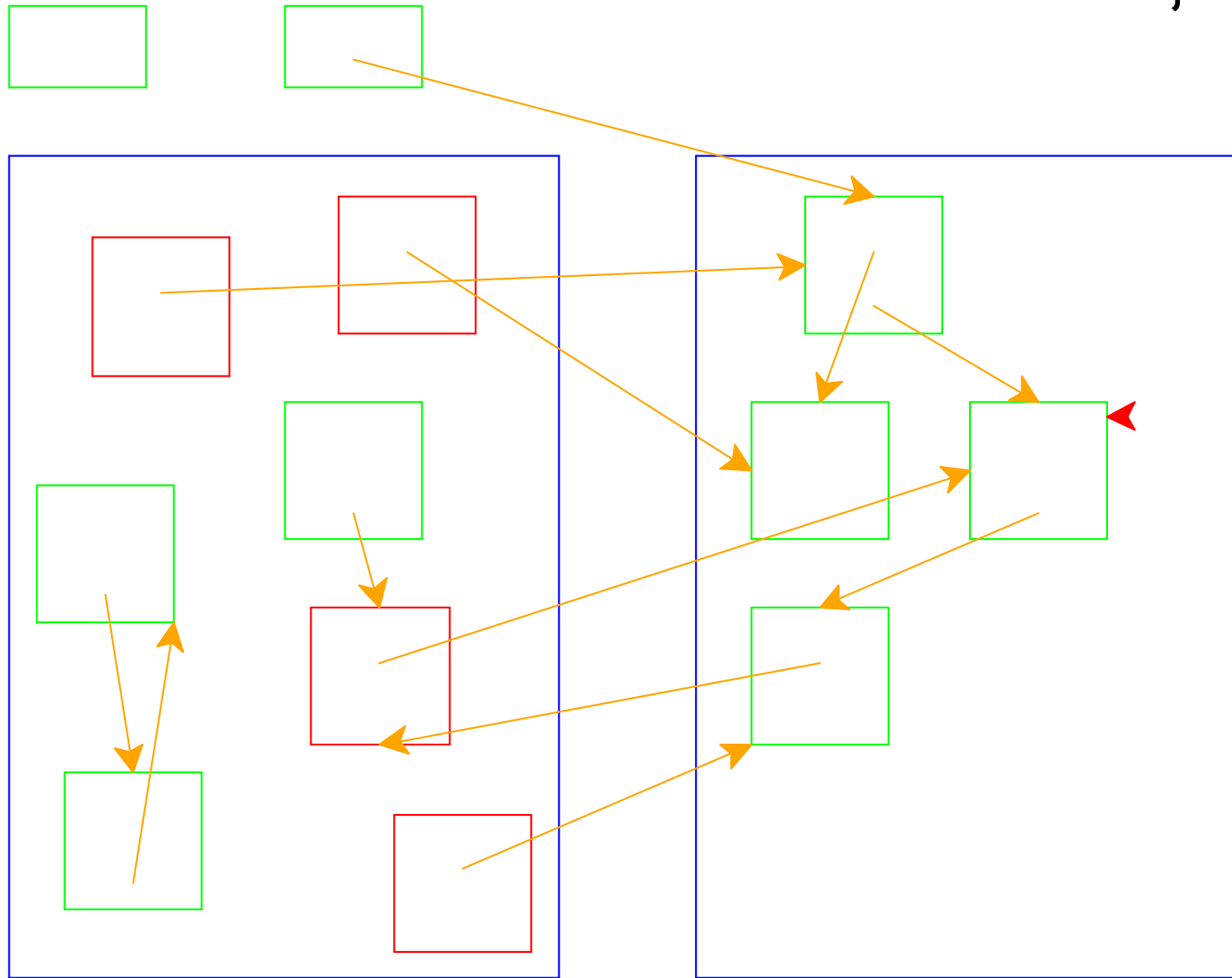
# Two-Space Collection

Nothing to color gray; increment the arrow

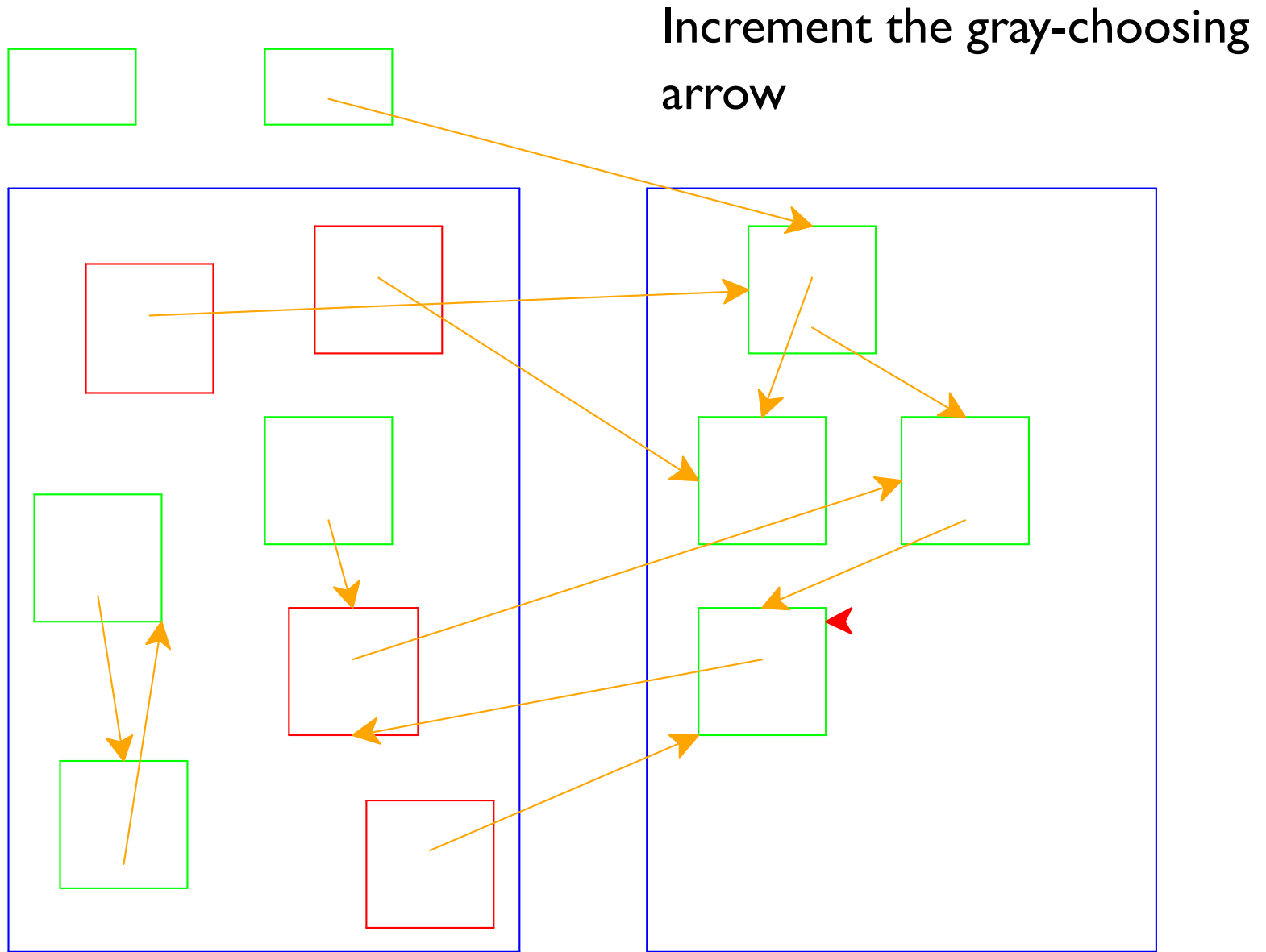


# Two-Space Collection

Color referenced object gray

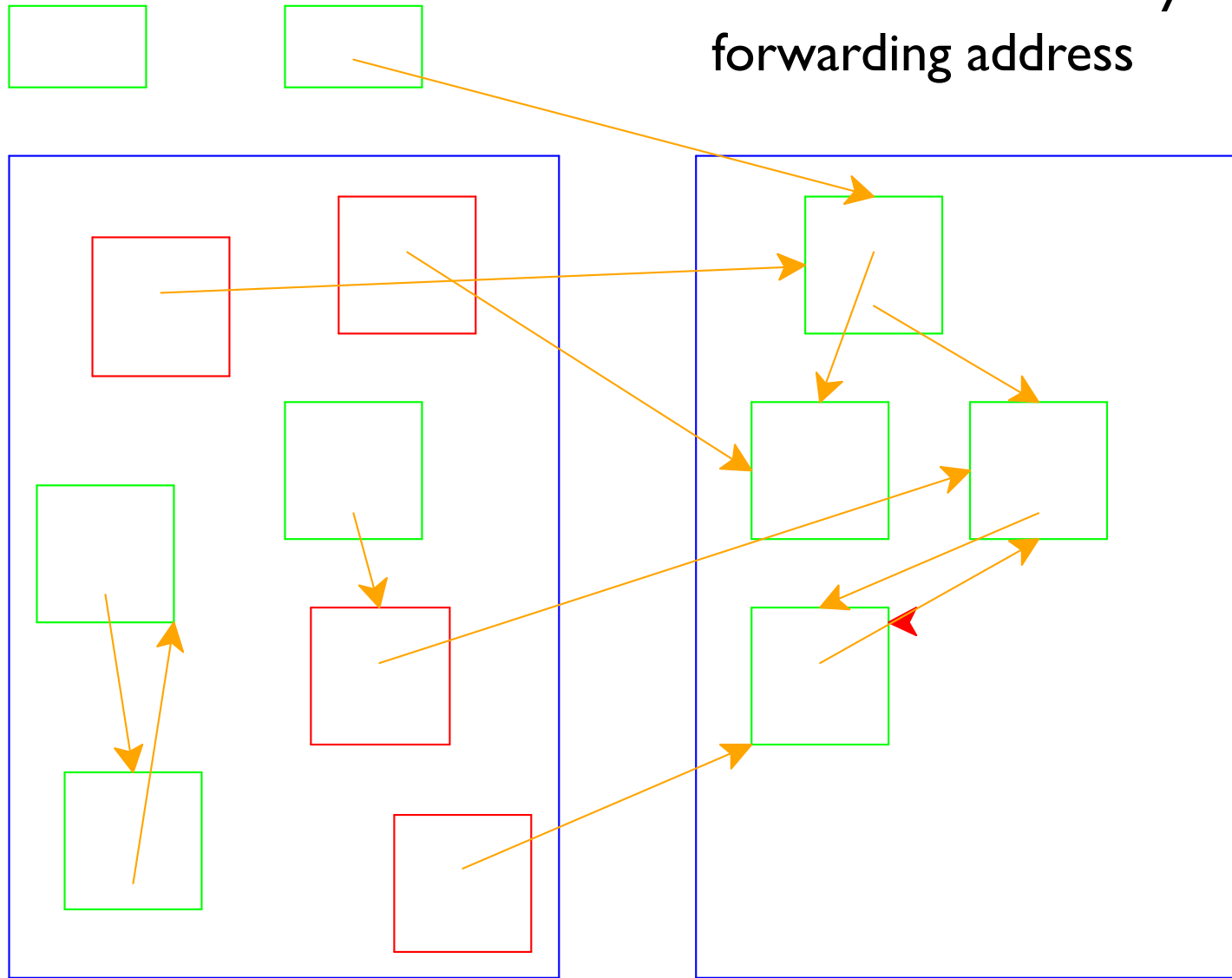


# Two-Space Collection



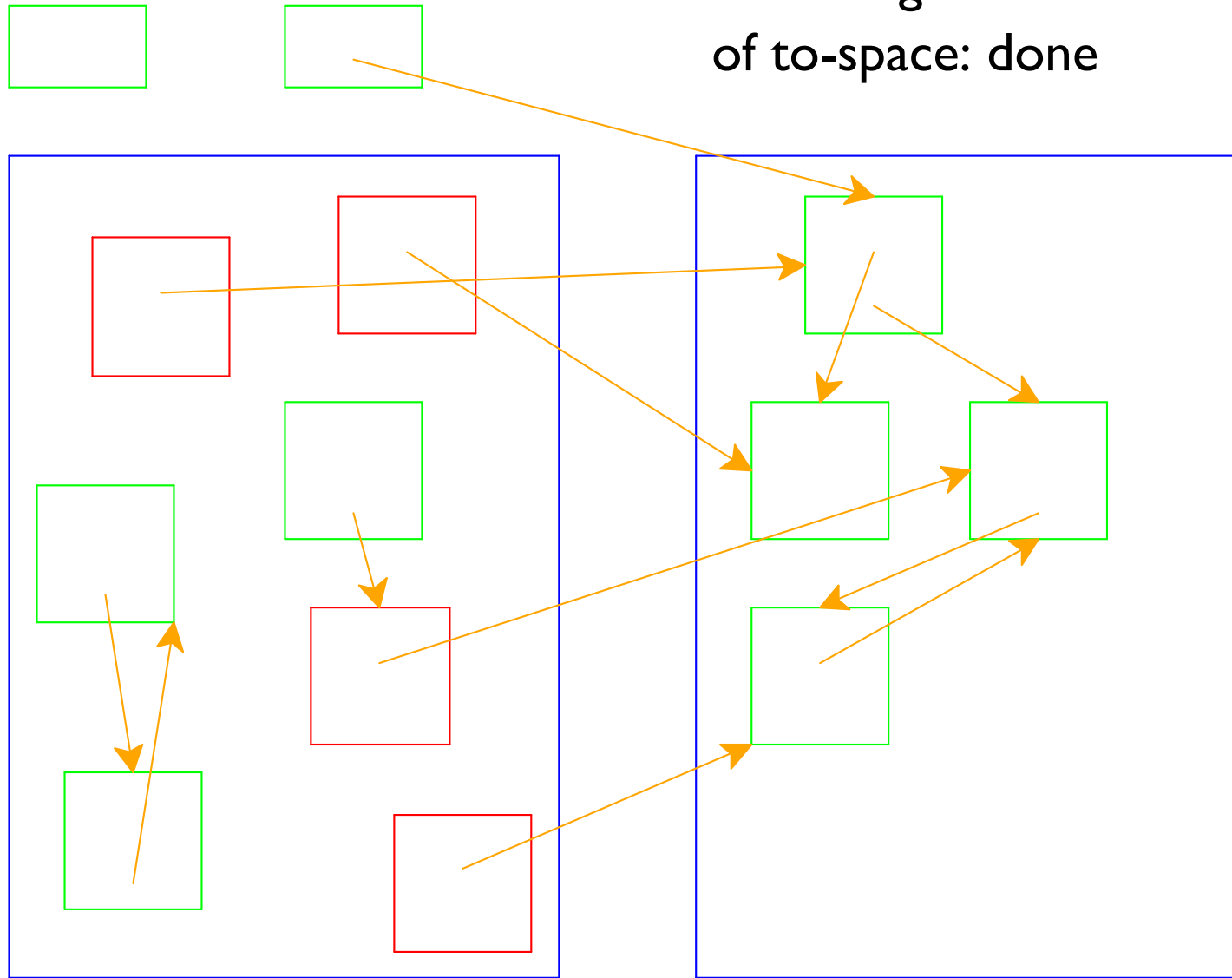
# Two-Space Collection

Referenced is already copied, use forwarding address



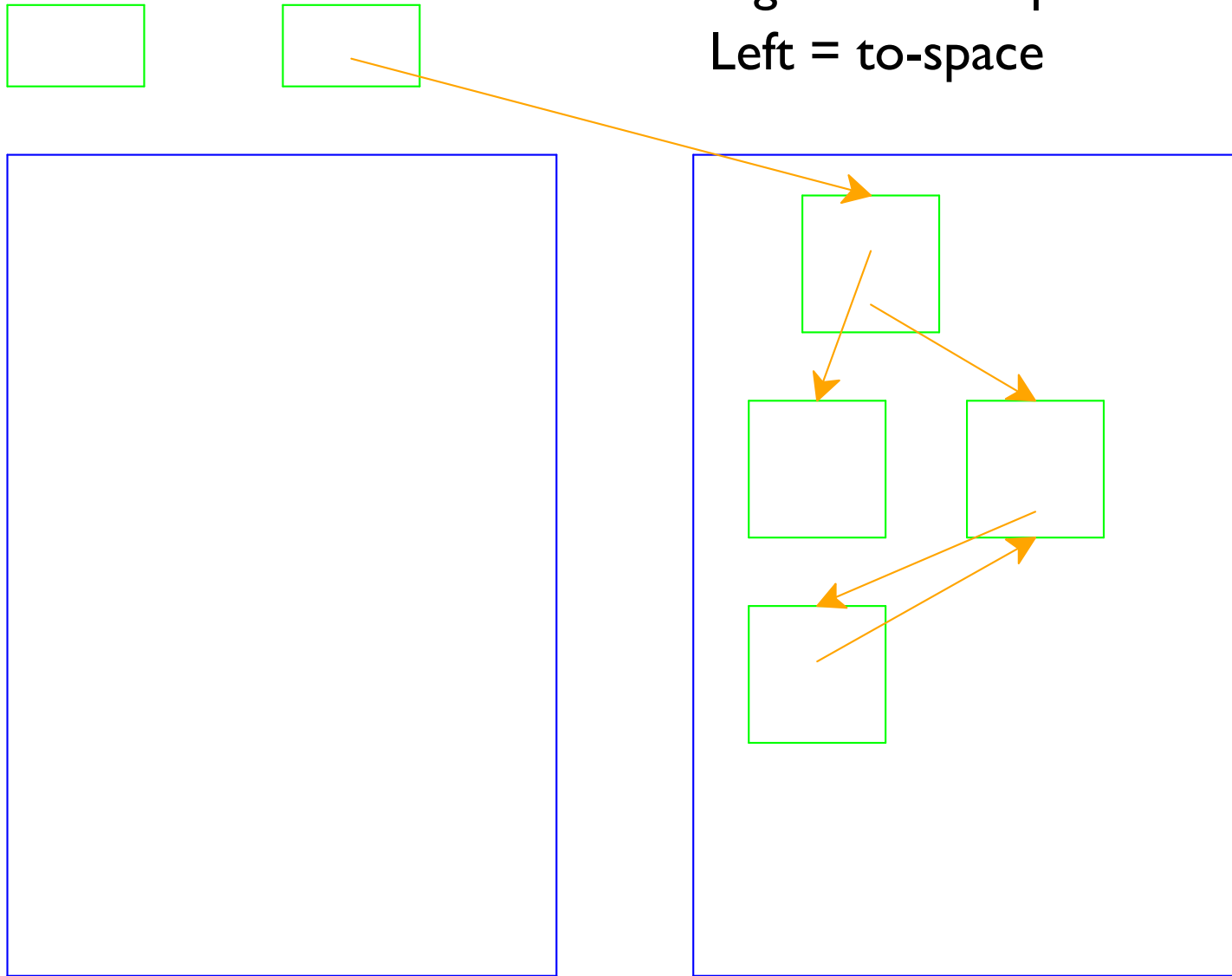
# Two-Space Collection

Choosing arrow reaches the end of to-space: done



# Two-Space Collection

Right = from-space  
Left = to-space



# Part 5



# Two-Space Collection on Vectors

- Everything is a number:
  - Some numbers are immediate integers
  - Some numbers are pointers
- An allocated object in memory starts with a tag, followed by a sequence of pointers and immediate integers
  - The tag describes the shape

# Two-Space Vector Example

- 26-byte memory (13 bytes for each space), 2 registers
  - Tag 1: one integer
  - Tag 2: one pointer
  - Tag 3: one integer, then one pointer

Register 1: 7

Register 2: 0

From: 1 75 2 0 3 2 10 3 2 2 3 1 4

# Two-Space Vector Example

- 26-byte memory (13 bytes for each space), 2 registers
  - Tag 1: one integer
  - Tag 2: one pointer
  - Tag 3: one integer, then one pointer

Register 1: 7

Register 2: 0

From:	1	75	2	0	3	2	10	3	2	2	3	1	4
Addr:	00	01	02	03	04	05	06	07	08	09	10	11	12

# Two-Space Vector Example

- 26-byte memory (13 bytes for each space), 2 registers
  - Tag 1: one integer
  - Tag 2: one pointer
  - Tag 3: one integer, then one pointer

	Register 1: 7						Register 2: 0						
From:	1	75	2	0	3	2	10	3	2	2	3	1	4
Addr:	00	01	02	03	04	05	06	07	08	09	10	11	12
	^		^		^			^			^		

# Two-Space Vector Example

- 26-byte memory (13 bytes for each space), 2 registers
  - Tag 1: one integer
  - Tag 2: one pointer
  - Tag 3: one integer, then one pointer

	Register 1: 7						Register 2: 0						
From:	1	75	2	0	3	2	10	3	2	2	3	1	4
Addr:	00	01	02	03	04	05	06	07	08	09	10	11	12
	^		^		^			^			^		
To:	0	0	0	0	0	0	0	0	0	0	0	0	0
	^												

# Two-Space Vector Example

- 26-byte memory (13 bytes for each space), 2 registers
  - Tag 1: one integer
  - Tag 2: one pointer
  - Tag 3: one integer, then one pointer

	Register 1: 0						Register 2: 0						
From:	1	75	2	0	3	2	10	99	0	2	3	1	4
Addr:	00	01	02	03	04	05	06	07	08	09	10	11	12
	^		^		^			^			^		
To:	3	2	2	0	0	0	0	0	0	0	0	0	0
	^												

# Two-Space Vector Example

- 26-byte memory (13 bytes for each space), 2 registers
  - Tag 1: one integer
  - Tag 2: one pointer
  - Tag 3: one integer, then one pointer

	Register 1: 0						Register 2: 3						
From:	99	3	2	0	3	2	10	99	0	2	3	1	4
Addr:	00	01	02	03	04	05	06	07	08	09	10	11	12
	^		^		^			^			^		
To:	3	2	2	1	75	0	0	0	0	0	0	0	0
	^												

# Two-Space Vector Example

- 26-byte memory (13 bytes for each space), 2 registers
  - Tag 1: one integer
  - Tag 2: one pointer
  - Tag 3: one integer, then one pointer

	Register 1: 0						Register 2: 3						
From:	99	3	99	5	3	2	10	99	0	2	3	1	4
Addr:	00	01	02	03	04	05	06	07	08	09	10	11	12
	^		^		^			^			^		
To:	3	2	5	1	75	2	0	0	0	0	0	0	0
				^									



# Two-Space Vector Example

- 26-byte memory (13 bytes for each space), 2 registers
  - Tag 1: one integer
  - Tag 2: one pointer
  - Tag 3: one integer, then one pointer

	Register 1: 0						Register 2: 3						
From:	99	3	99	5	3	2	10	99	0	2	3	1	4
Addr:	00	01	02	03	04	05	06	07	08	09	10	11	12
	^		^		^			^			^		
To:	3	2	5	1	75	2	0	0	0	0	0	0	0
						^							

# Two-Space Vector Example

- 26-byte memory (13 bytes for each space), 2 registers
  - Tag 1: one integer
  - Tag 2: one pointer
  - Tag 3: one integer, then one pointer

	Register 1: 0						Register 2: 3						
From:	99	3	99	5	3	2	10	99	0	2	3	1	4
Addr:	00	01	02	03	04	05	06	07	08	09	10	11	12
	^		^		^			^			^		
To:	3	2	5	1	75	2	3	0	0	0	0	0	0
								^					